Adapting TCP for Reconfigurable Datacenter Networks

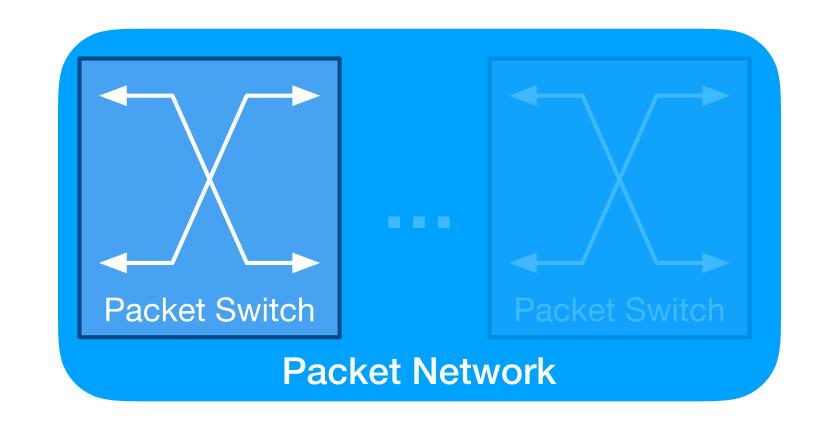
Matthew K. Mukerjee*+, Christopher Canel*, Weiyang Wang°, Daehyeok Kim*+,

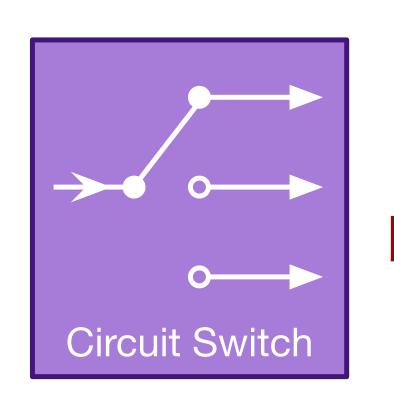
Srinivasan Seshan*, Alex C. Snoeren°

*Carnegie Mellon University, °UC San Diego, †Nefeli Networks, ‡Microsoft Research

February 26, 2020

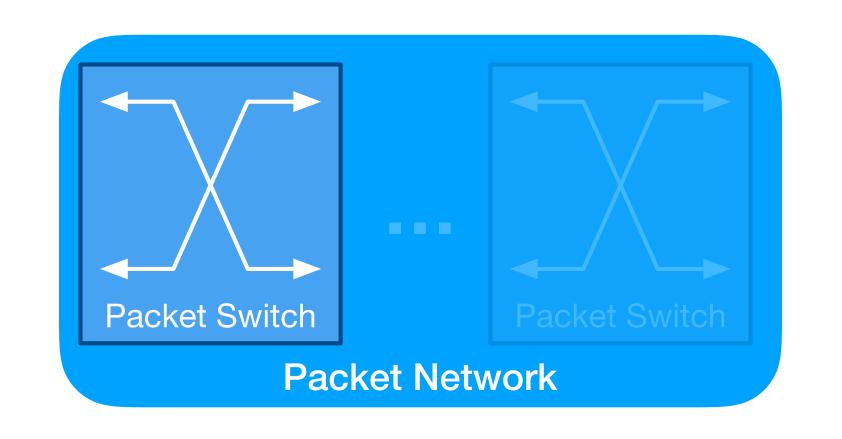
all-to-all connectivity

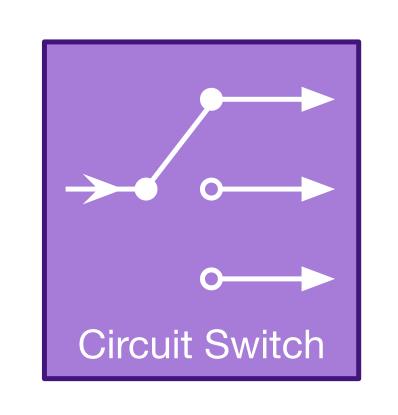




higher bandwidth, between certain racks

all-to-all connectivity





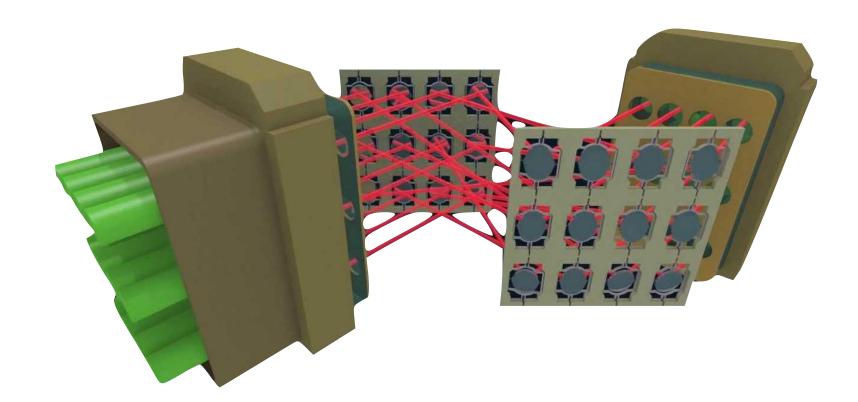
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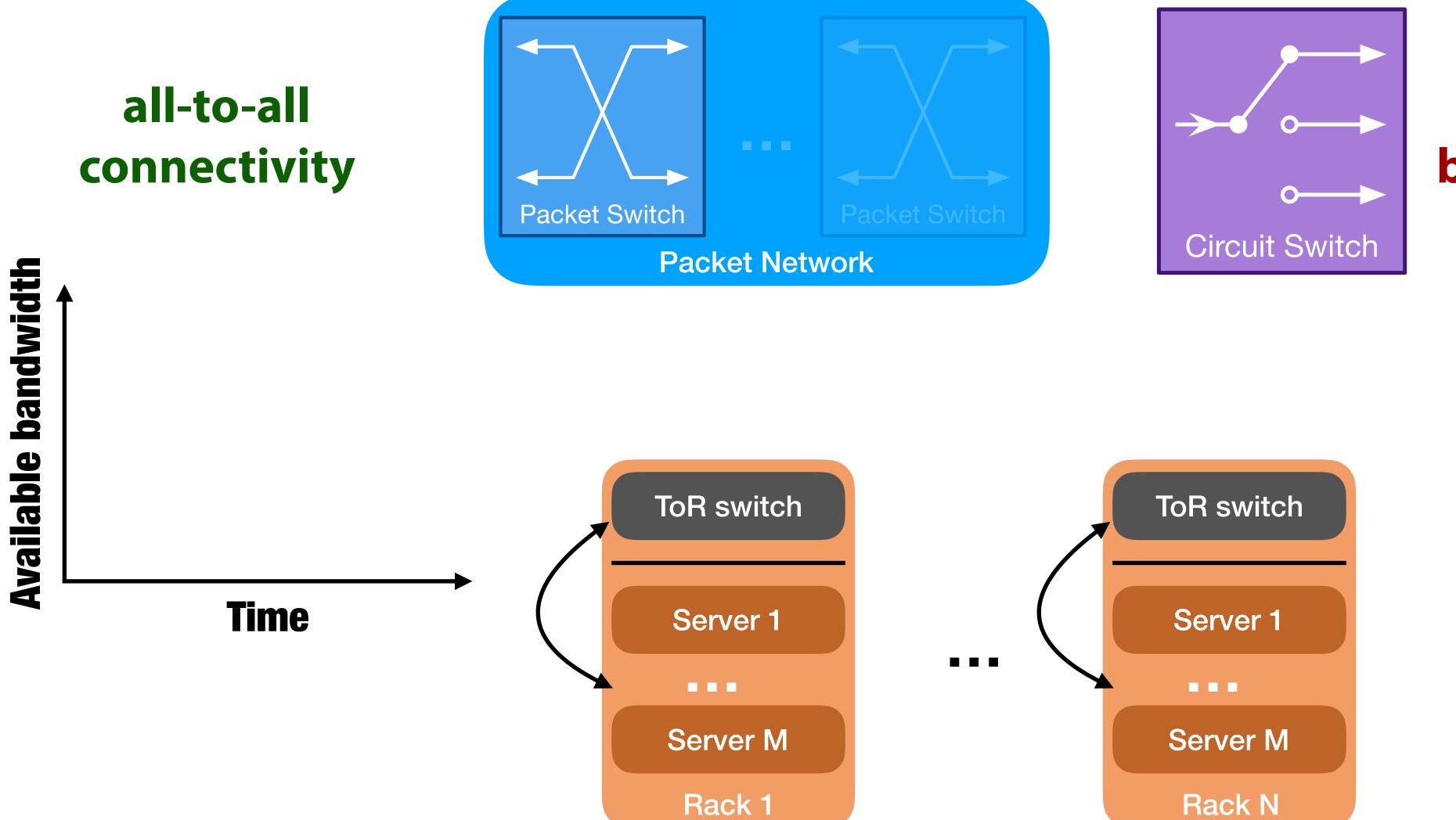
60GHz wireless



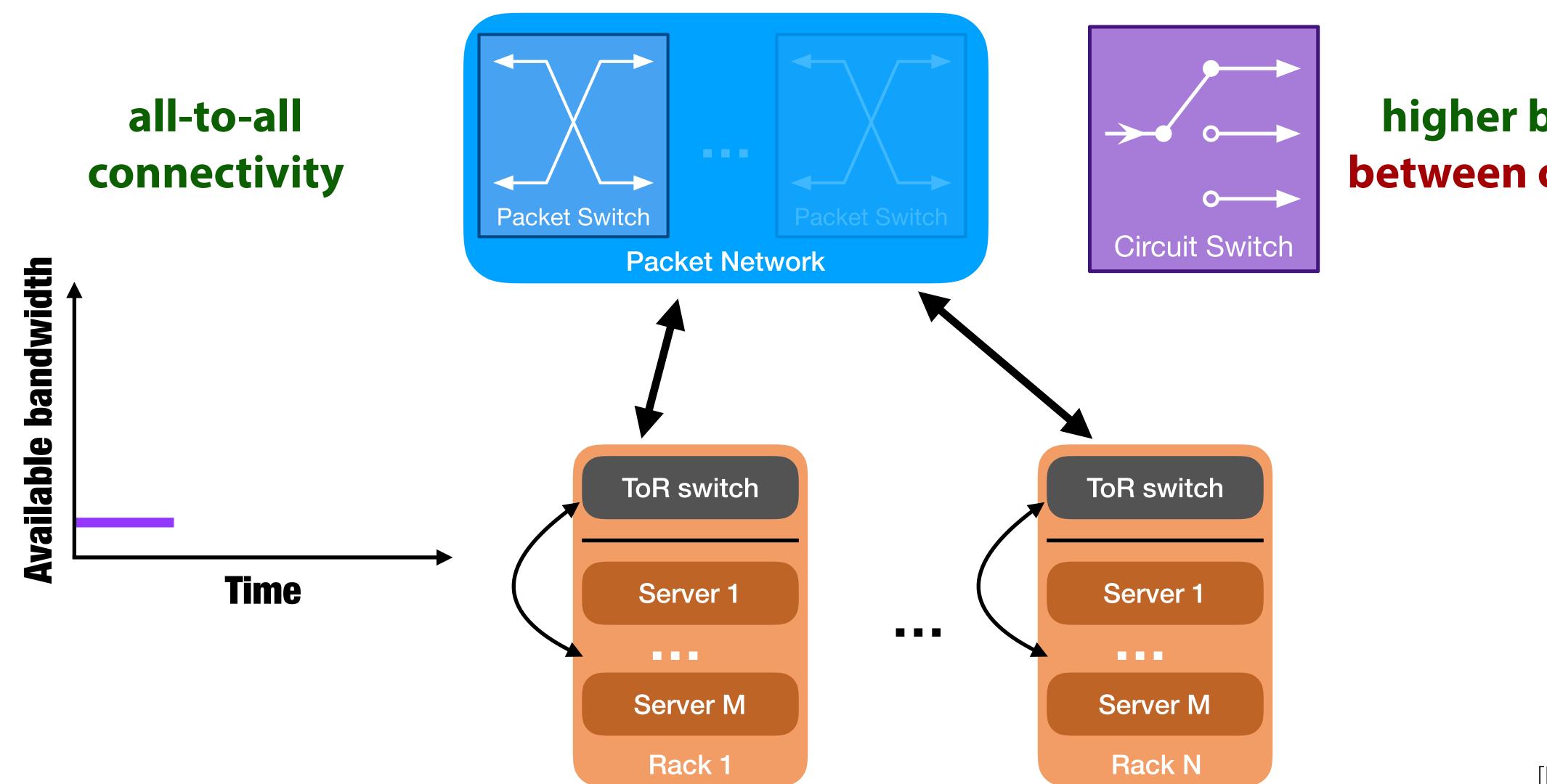
free-space optics



optical circuits

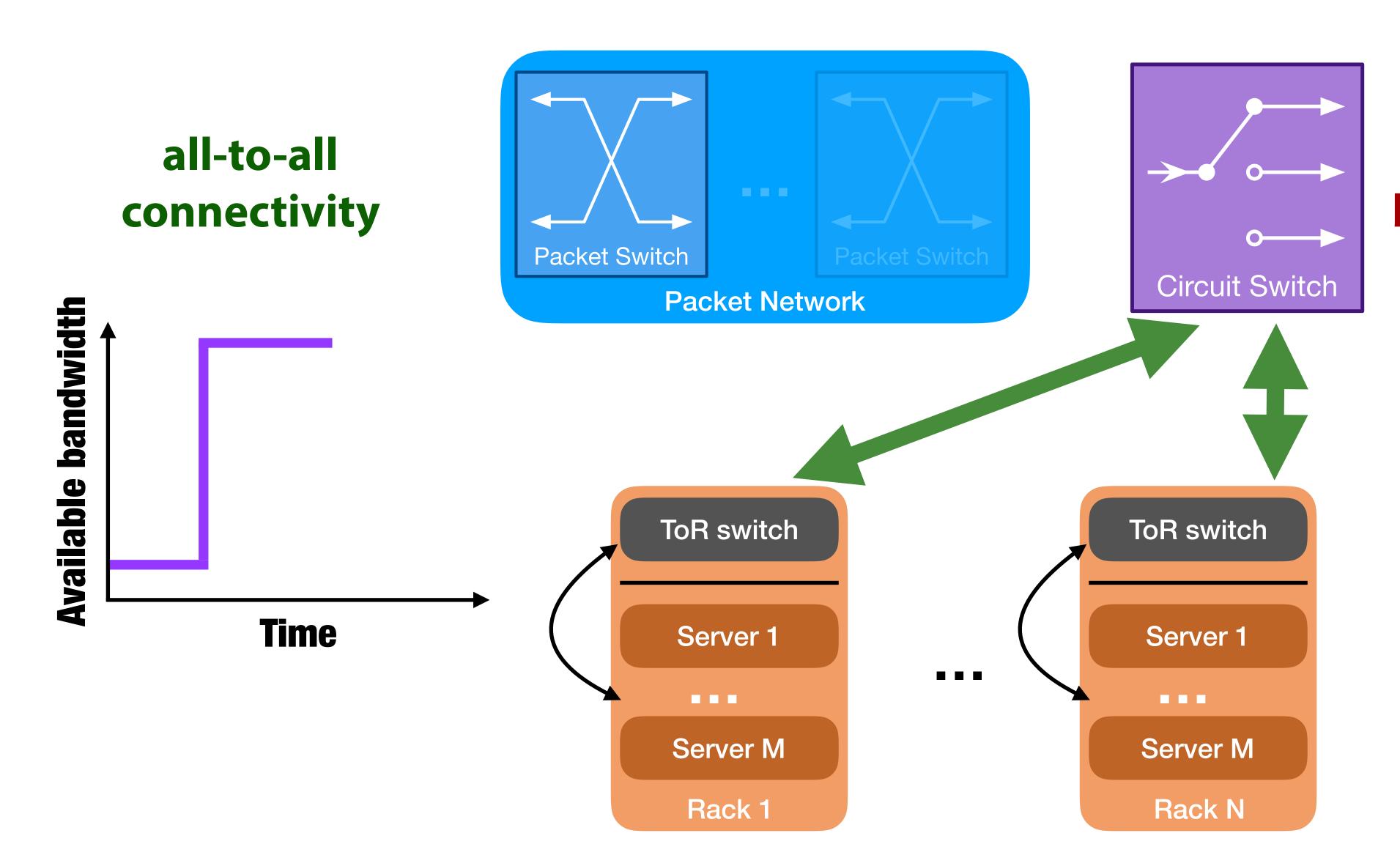


higher bandwidth, between certain racks

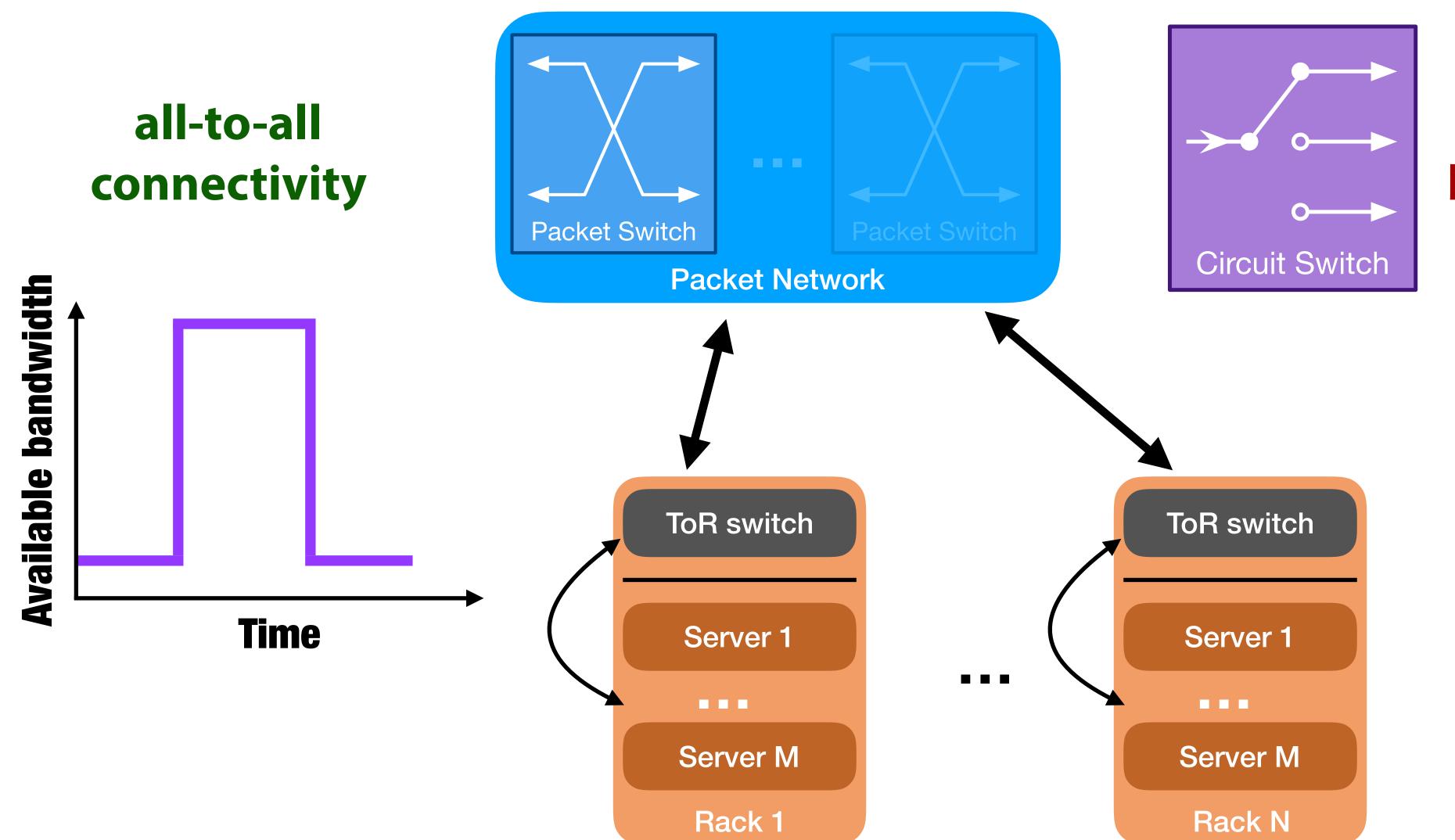


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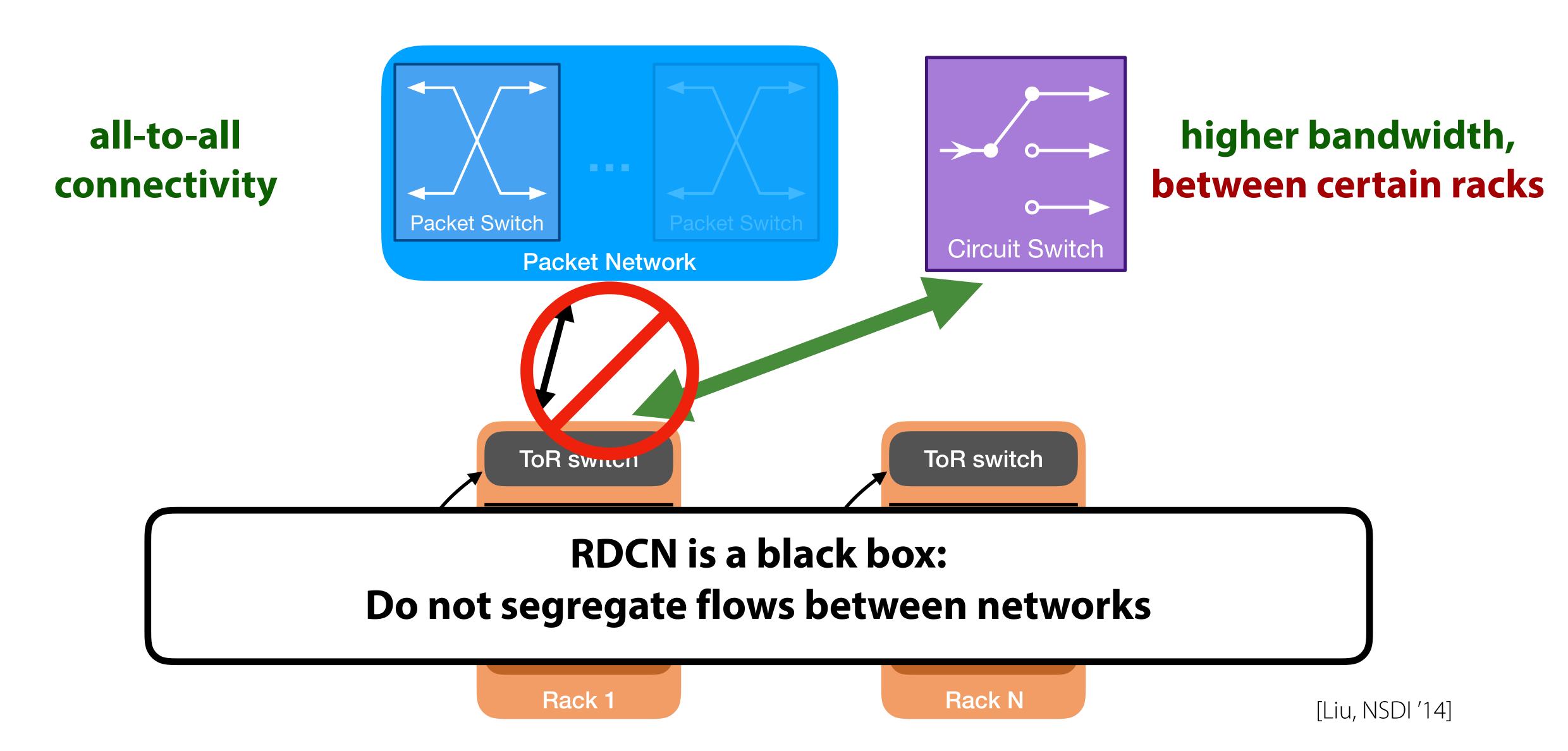
[Liu, NSDI '14]



higher bandwidth, between certain racks



higher bandwidth, between certain racks



2010: RDCNs speed up DC workloads

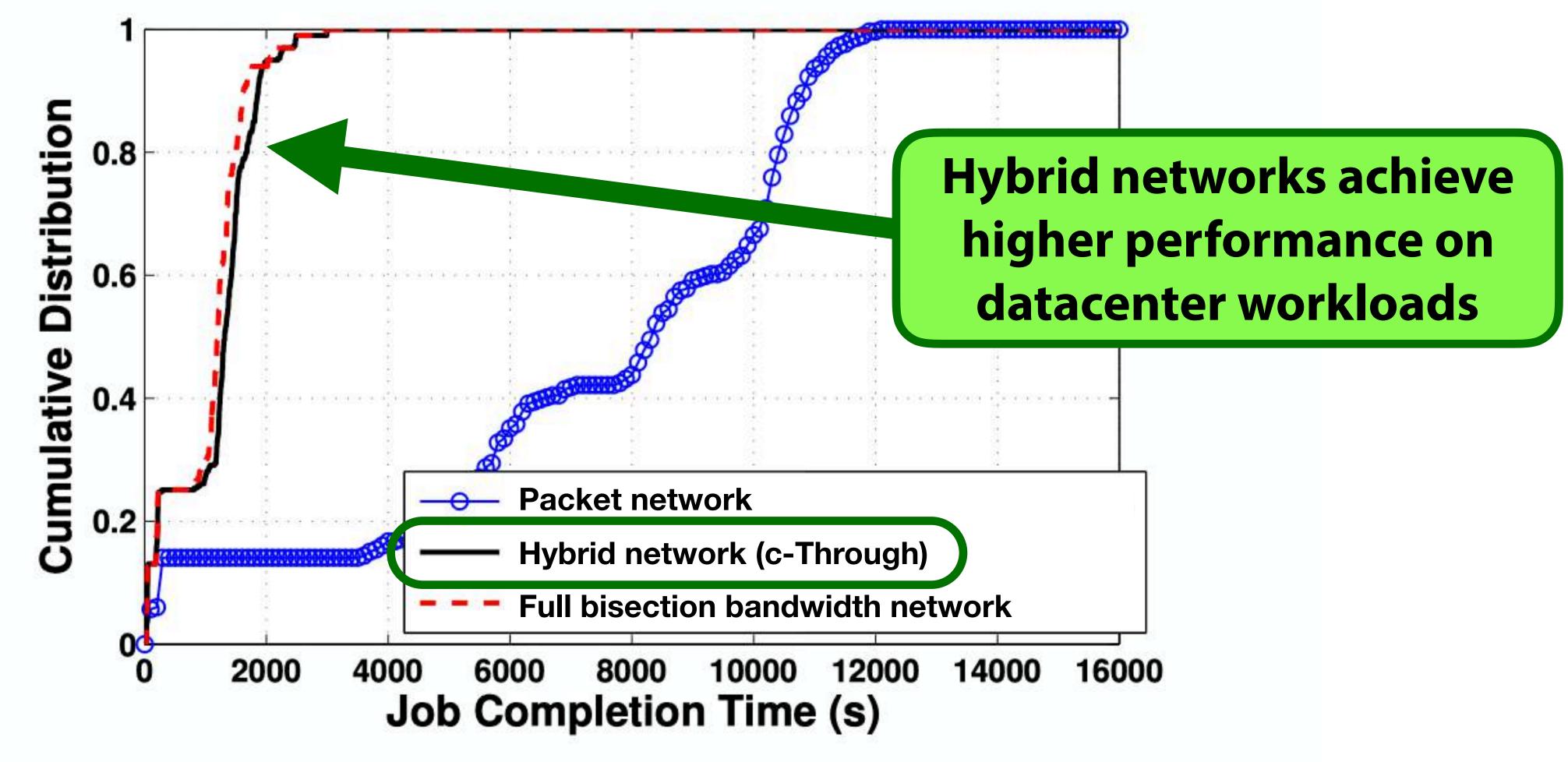
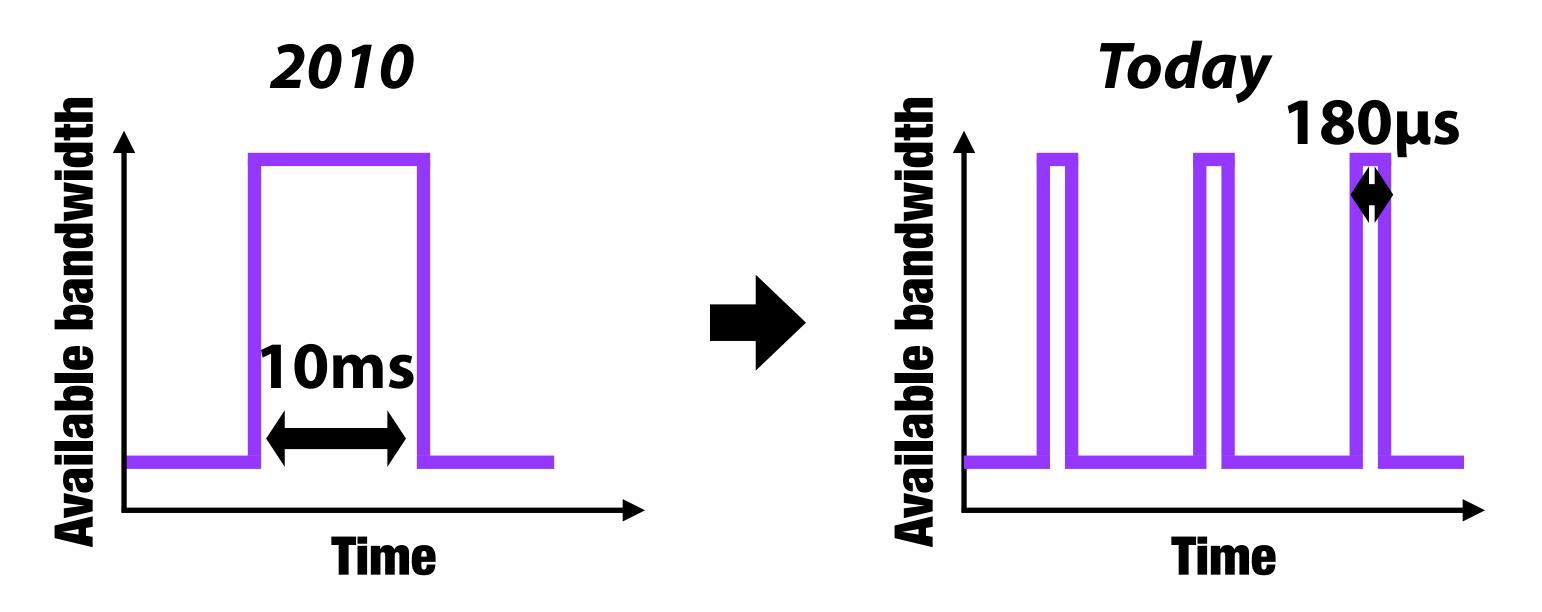


Figure 9: The completion of Hadoop Gridmix tasks

Today's RDCNs reconfigure 10x as often

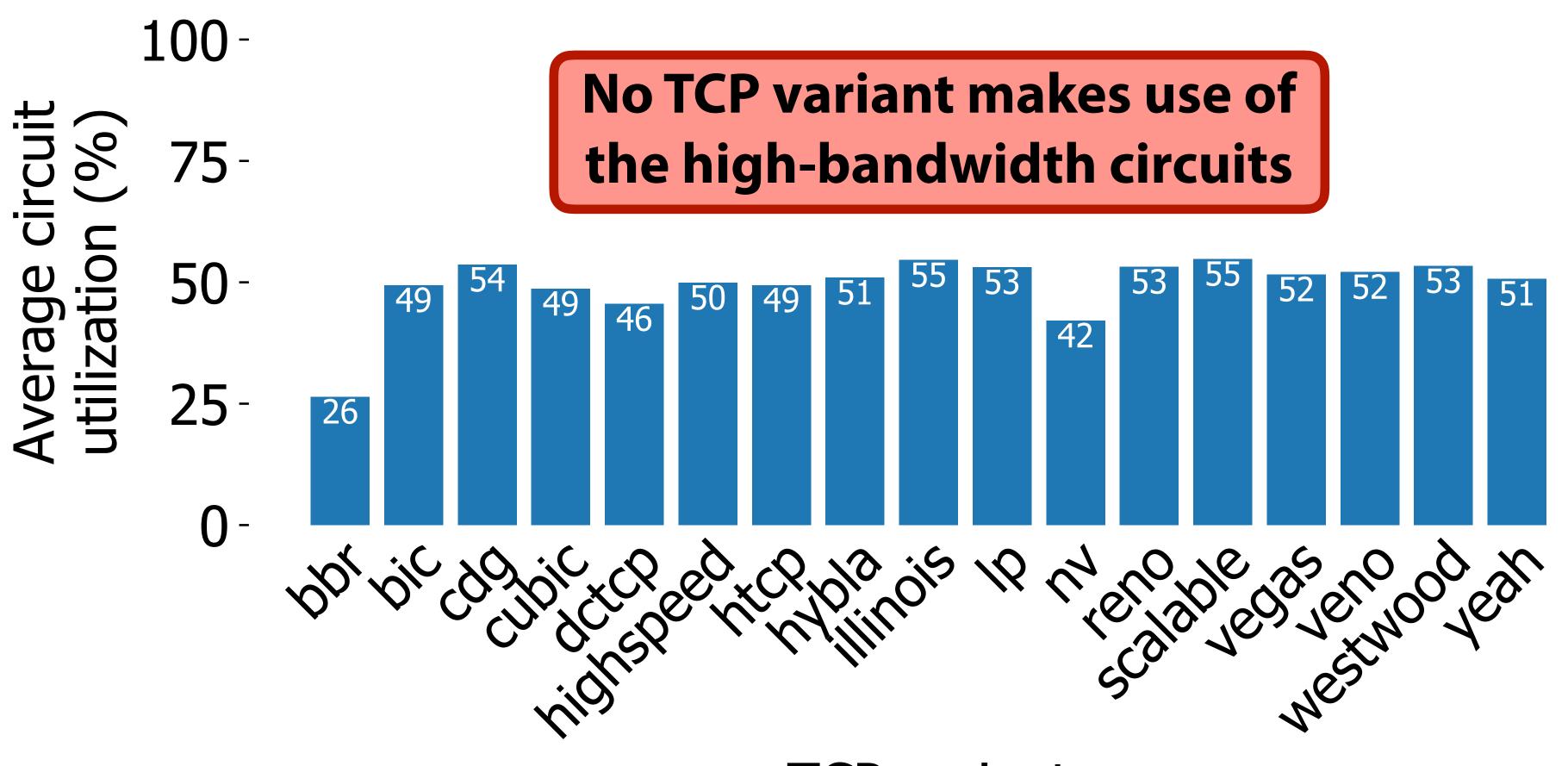
Advances in circuit switch technology have led to a 10x reduction in reconfiguration delay \Rightarrow today, circuits can reconfigure much more frequently



Better for datacenters: More flexibility to support dynamic workloads **Better for hosts:** Less data must be available to saturate higher bandwidth NW

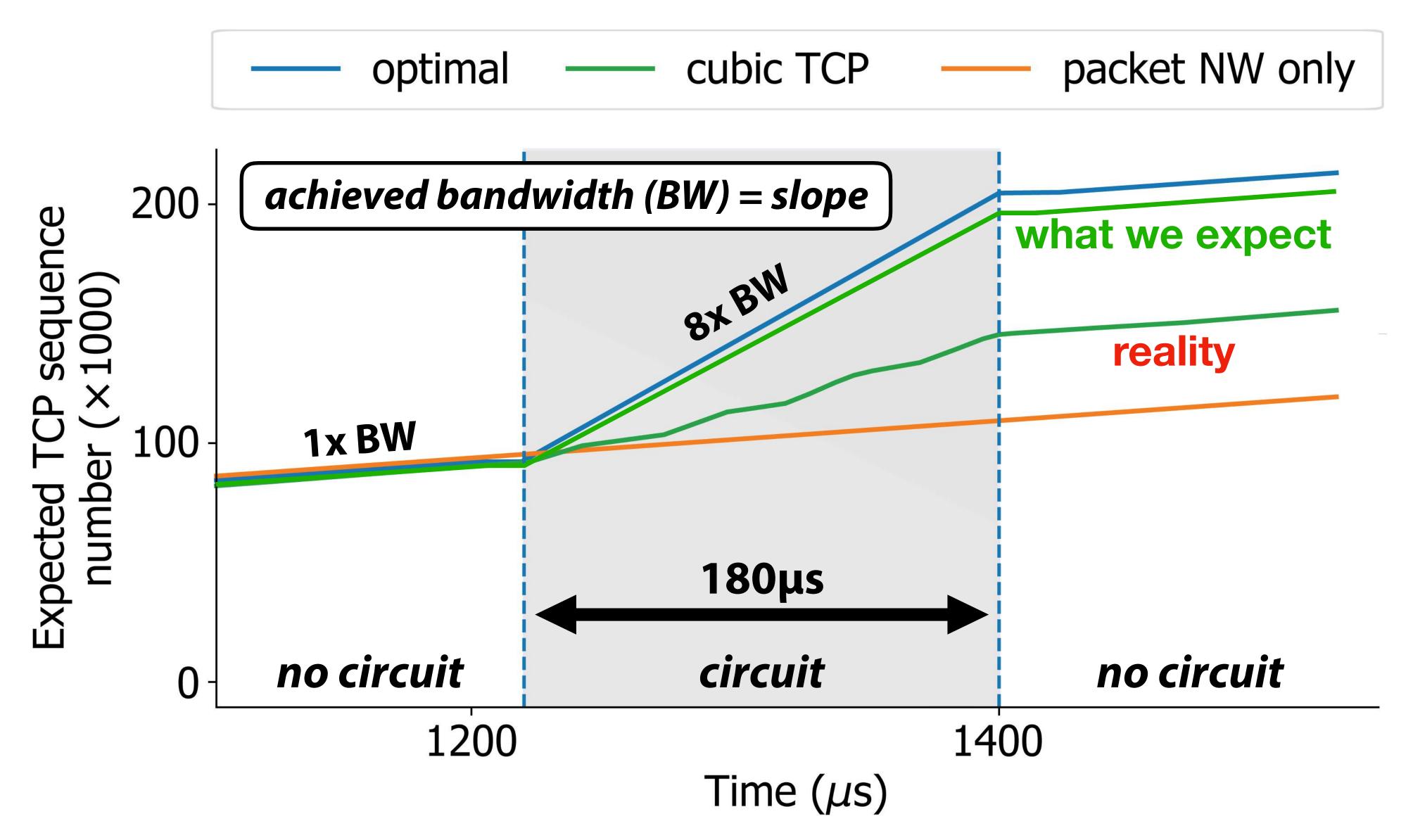
Short-lived circuits pose a problem for TCP

16 flows from rack 1 to rack 2; packet network: 10 Gb/s; circuit network: 80 Gb/s



TCP variant

TCP cannot ramp up during short circuits



What is the problem?

All TCP variants are designed to adapt to changing network conditions

E.g., congestion, bottleneck links, RTT

But **bandwidth fluctuations** in modern RDCN are an order of magnitude **more frequent** (10x shorter circuit duration) and **more substantial** (10x higher bandwidth) than TCP is designed to handle

• RDCNs break the implicit assumption of relatively-stable network conditions

This requires an order-of-magnitude shift in how fast TCP reacts

This talk: Our 2-part solution

In-network: Use information about upcoming circuits to transparently "trick" TCP into ramping up more aggressively

High utilization, at the cost of tail latency

At endhosts: New TCP variant, reTCP, that explicitly reacts to circuit state changes

Mitigates tail latency penalty

The two techniques can be deployed separately, but work best together

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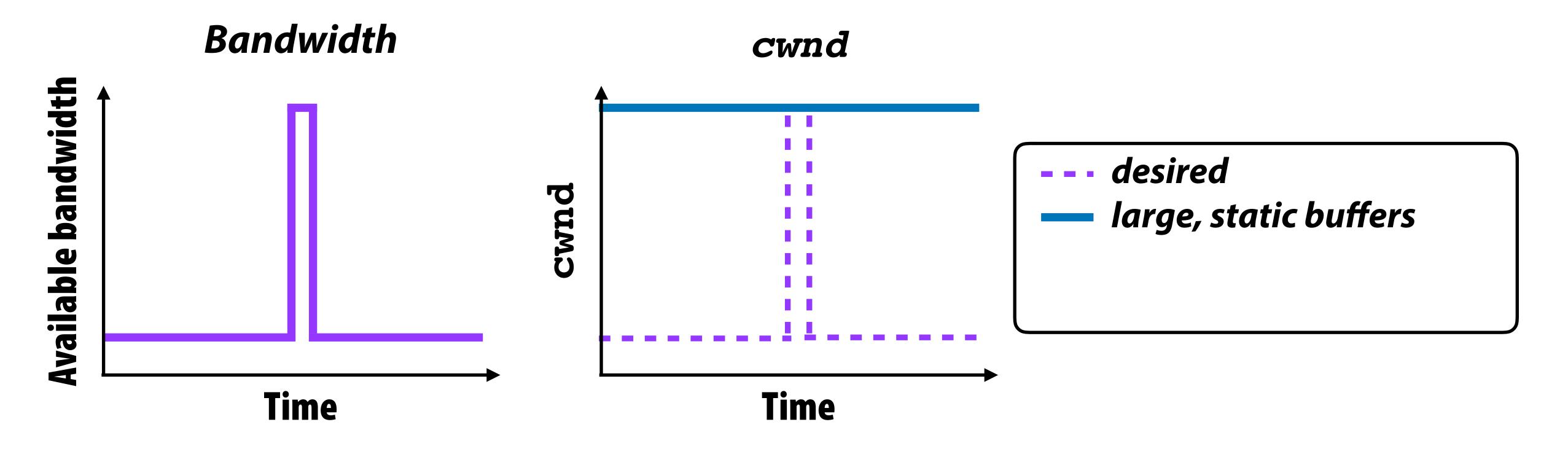
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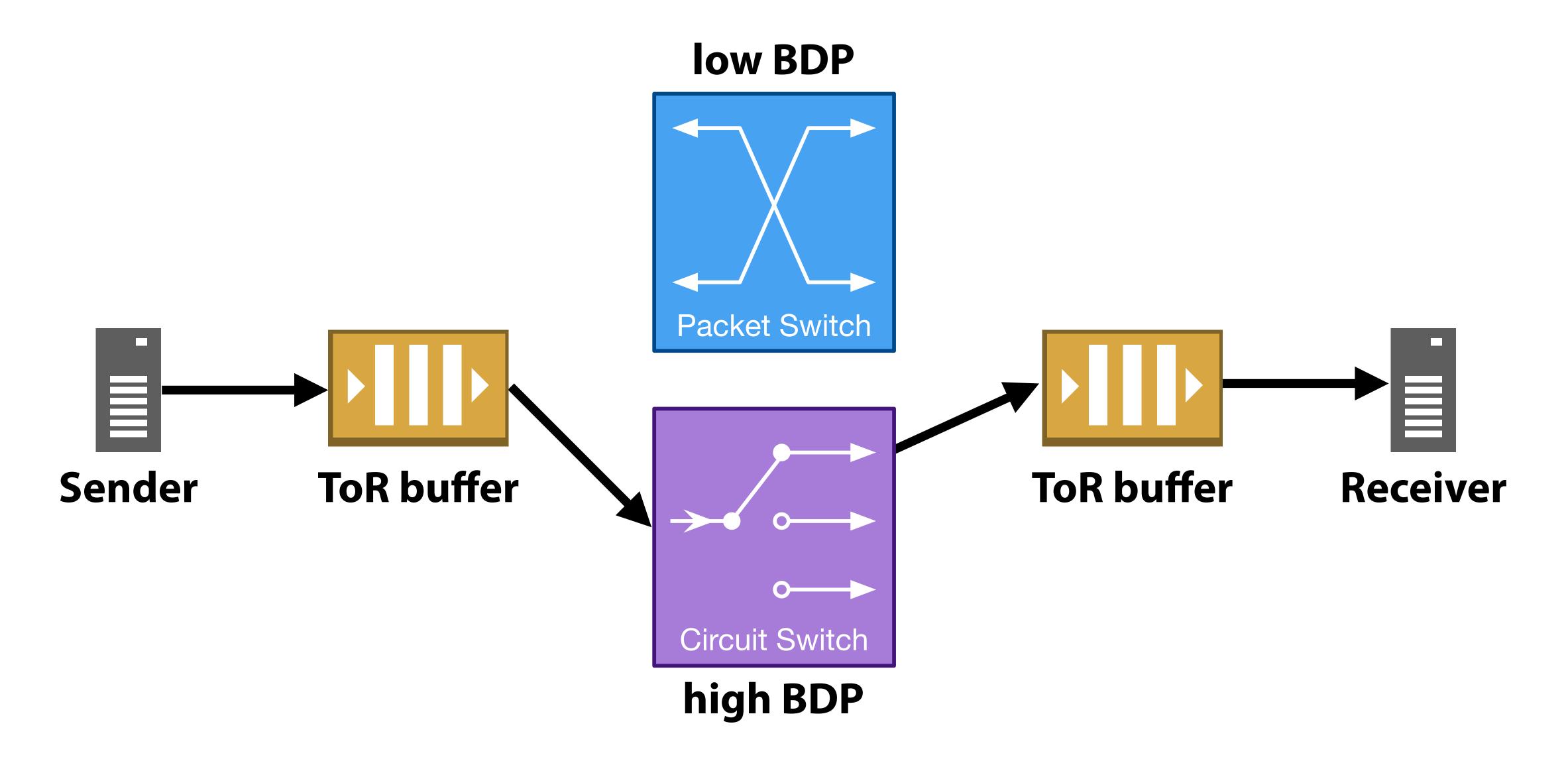
Mitigates tail latency penalty

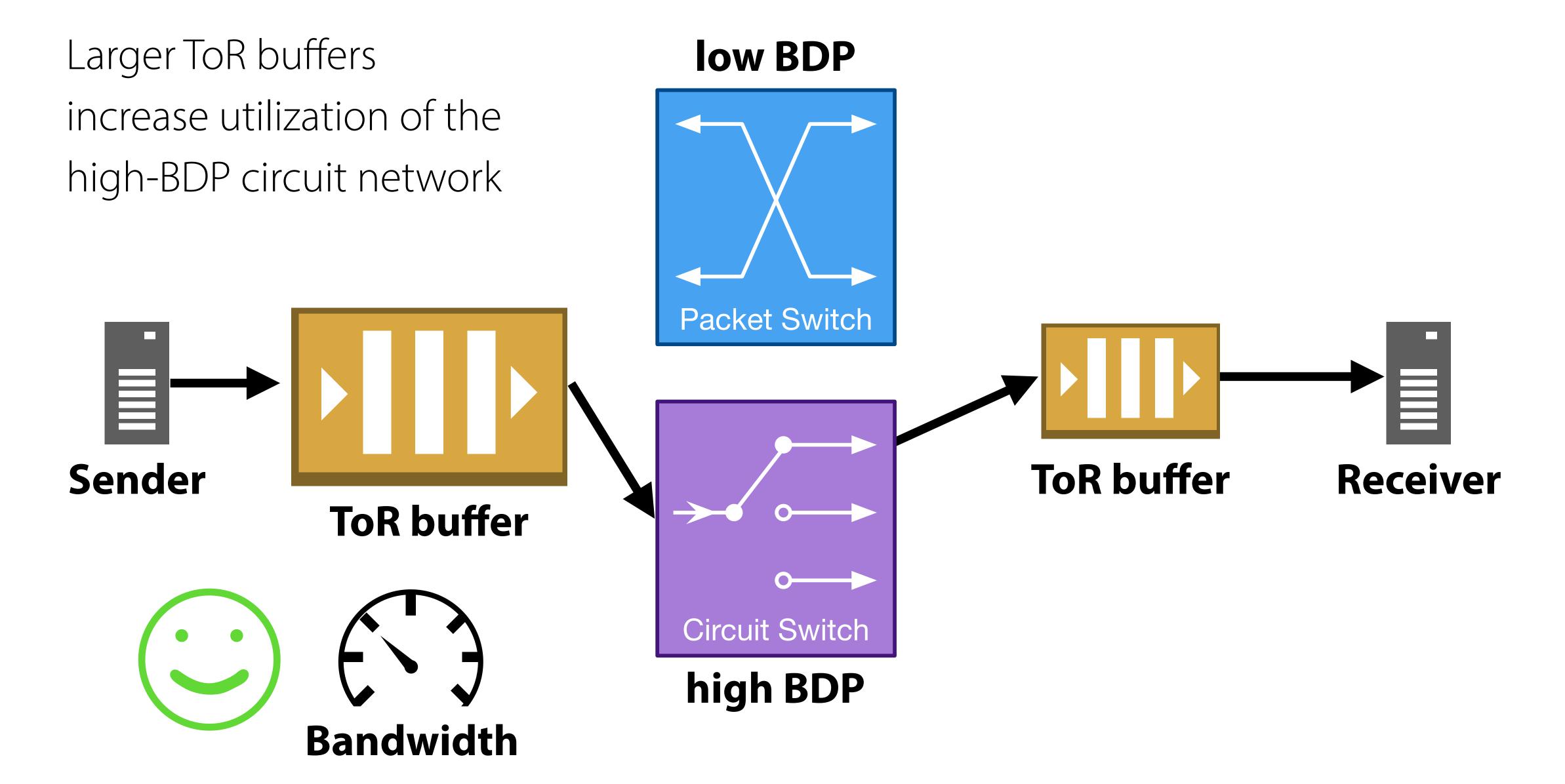
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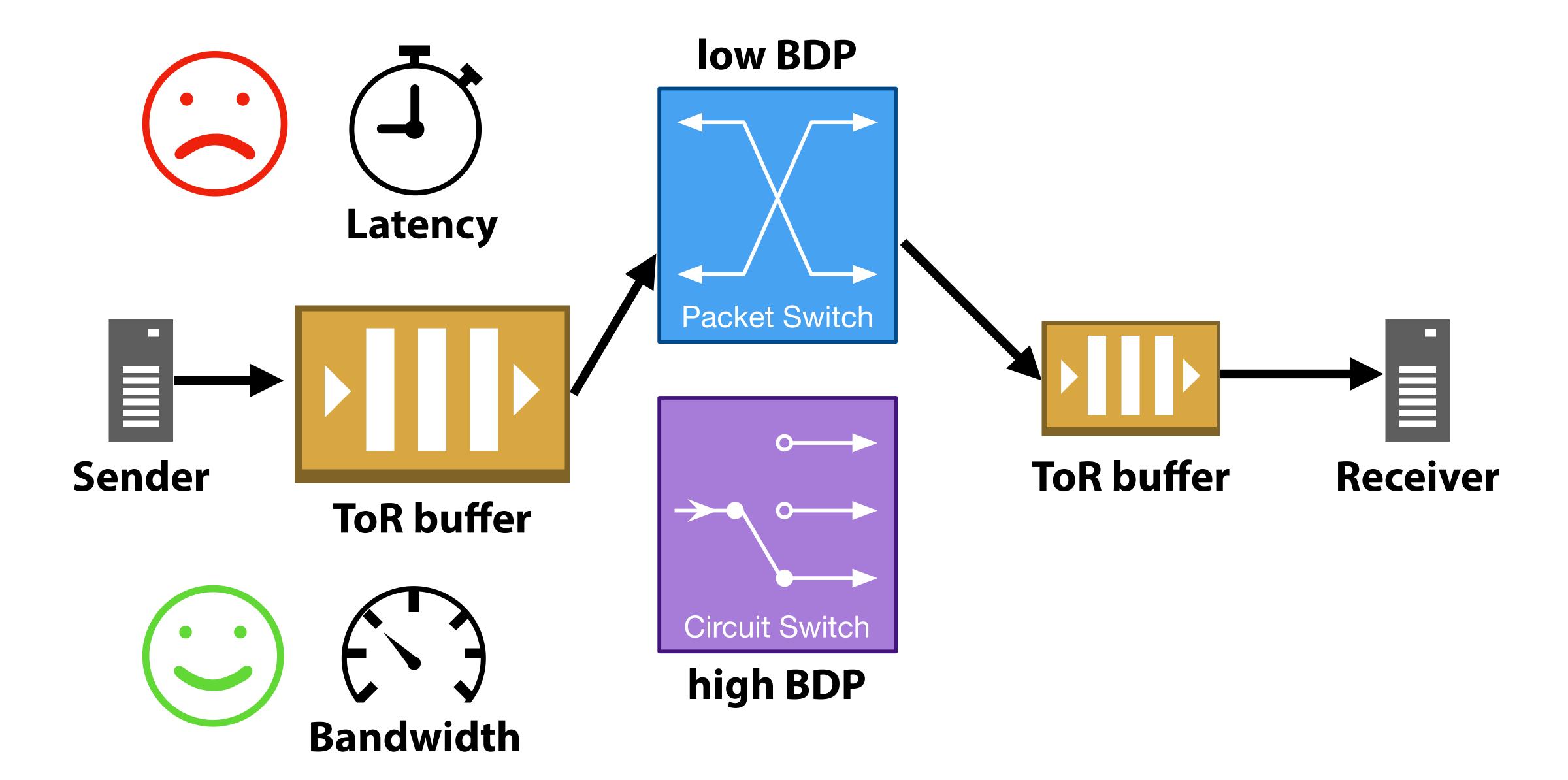
Want we want: TCP's congestion window (cwnd) to parallel the BW fluctuations

First attempt: Make cwnd large all the time How? Use large ToR buffers

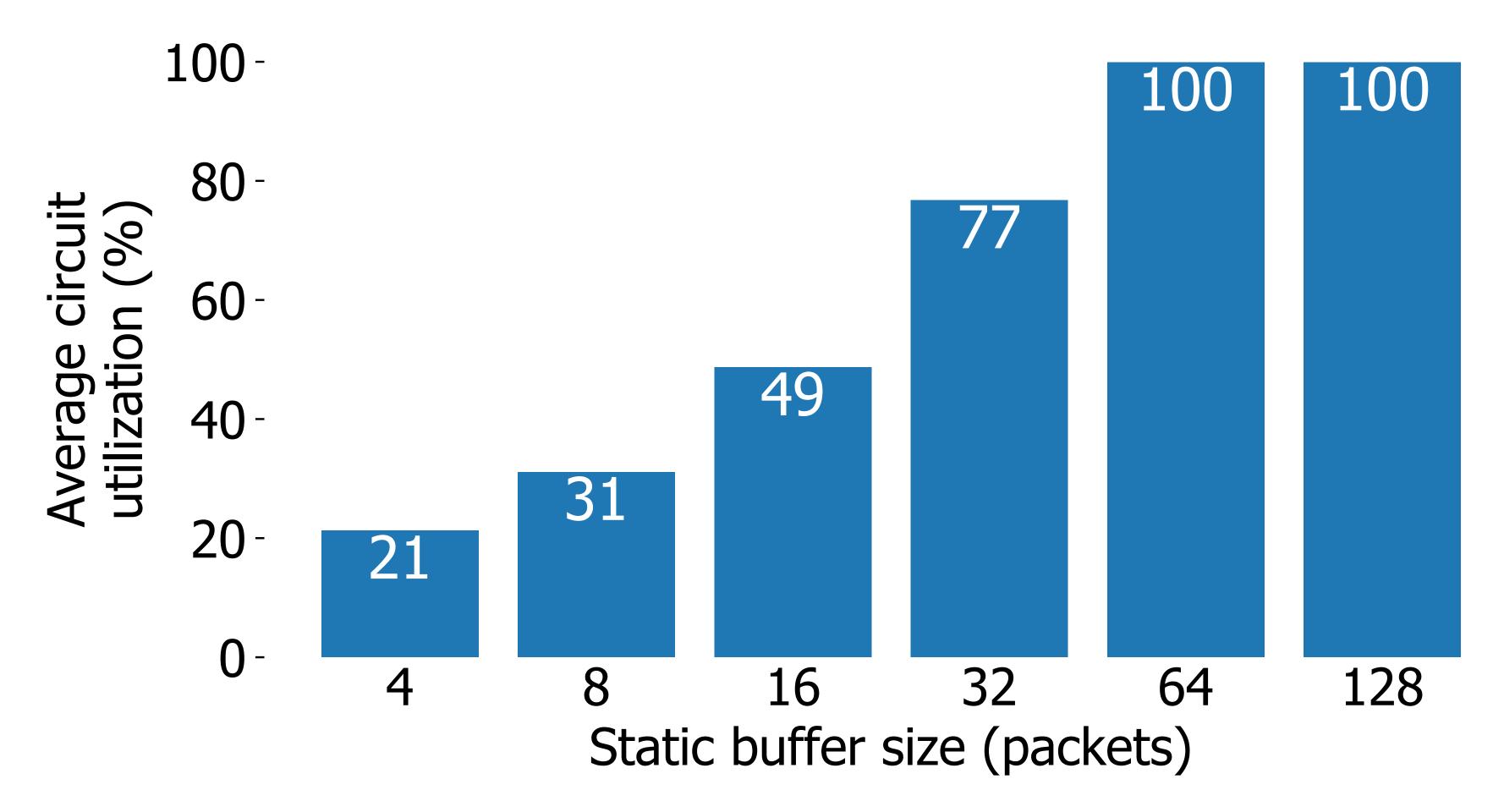






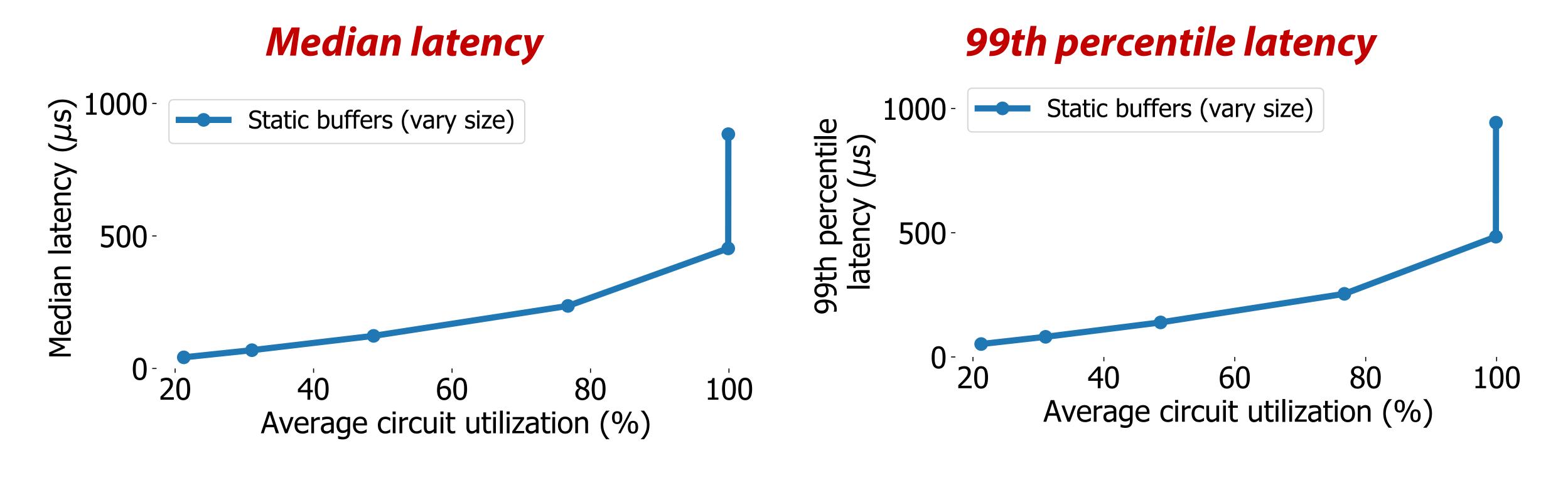


Large queues increase utilization...



16 flows from rack 1 to rack 2; packet network: 10 Gb/s; circuit network: 80 Gb/s

...but result in high latency

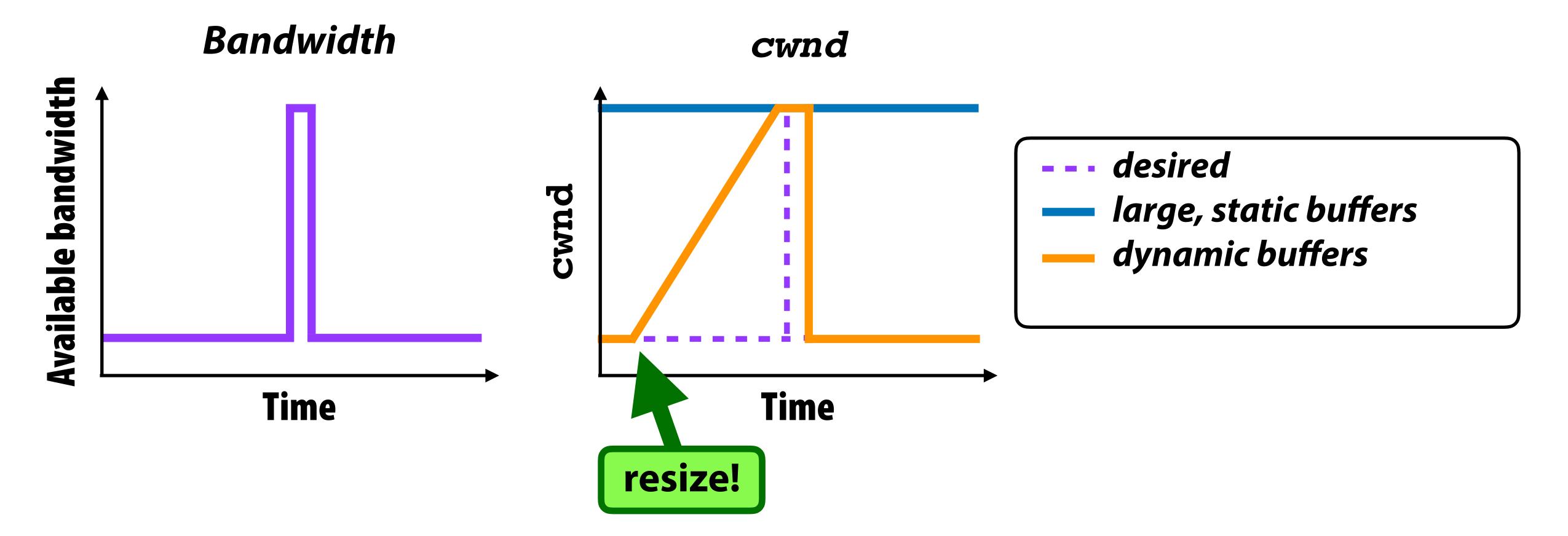


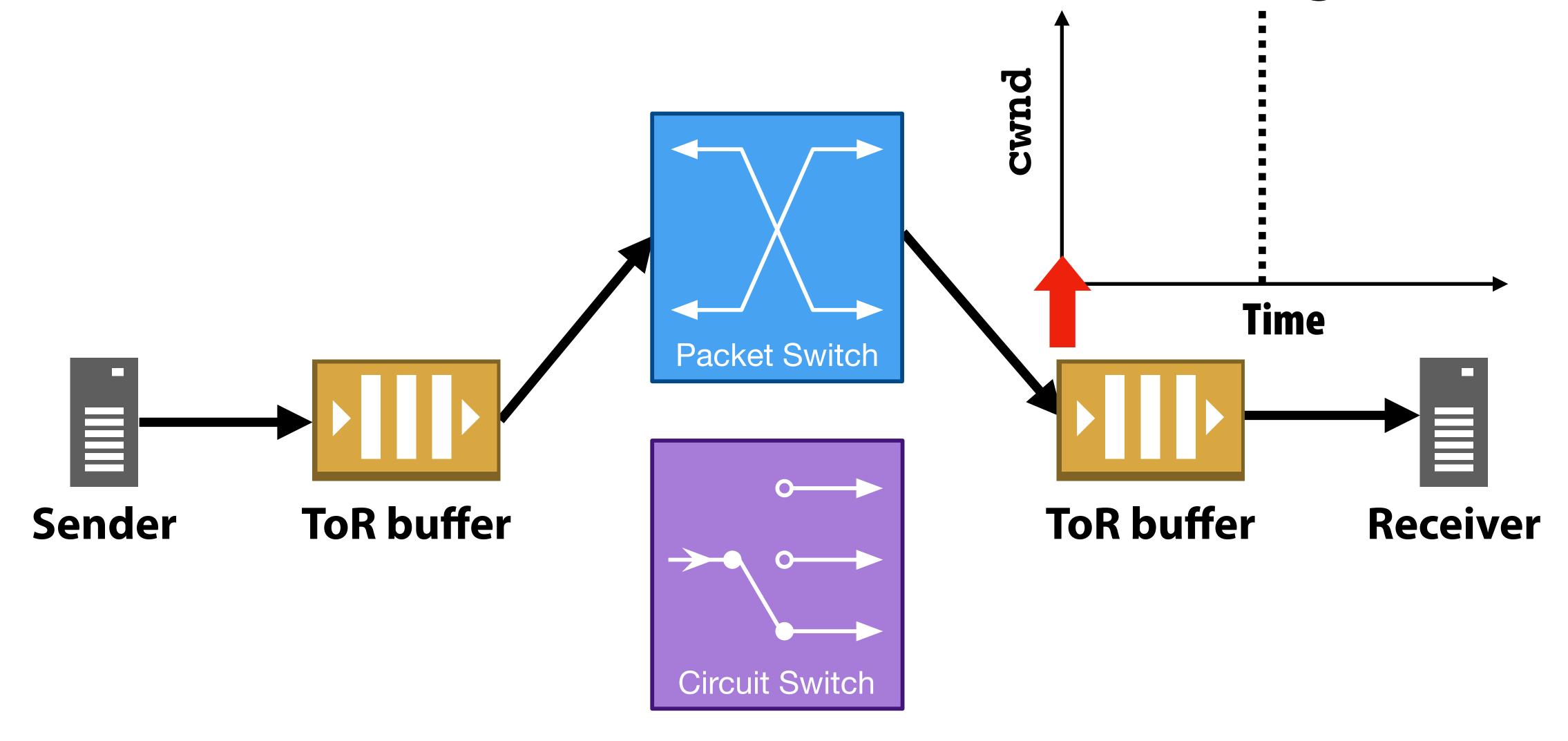
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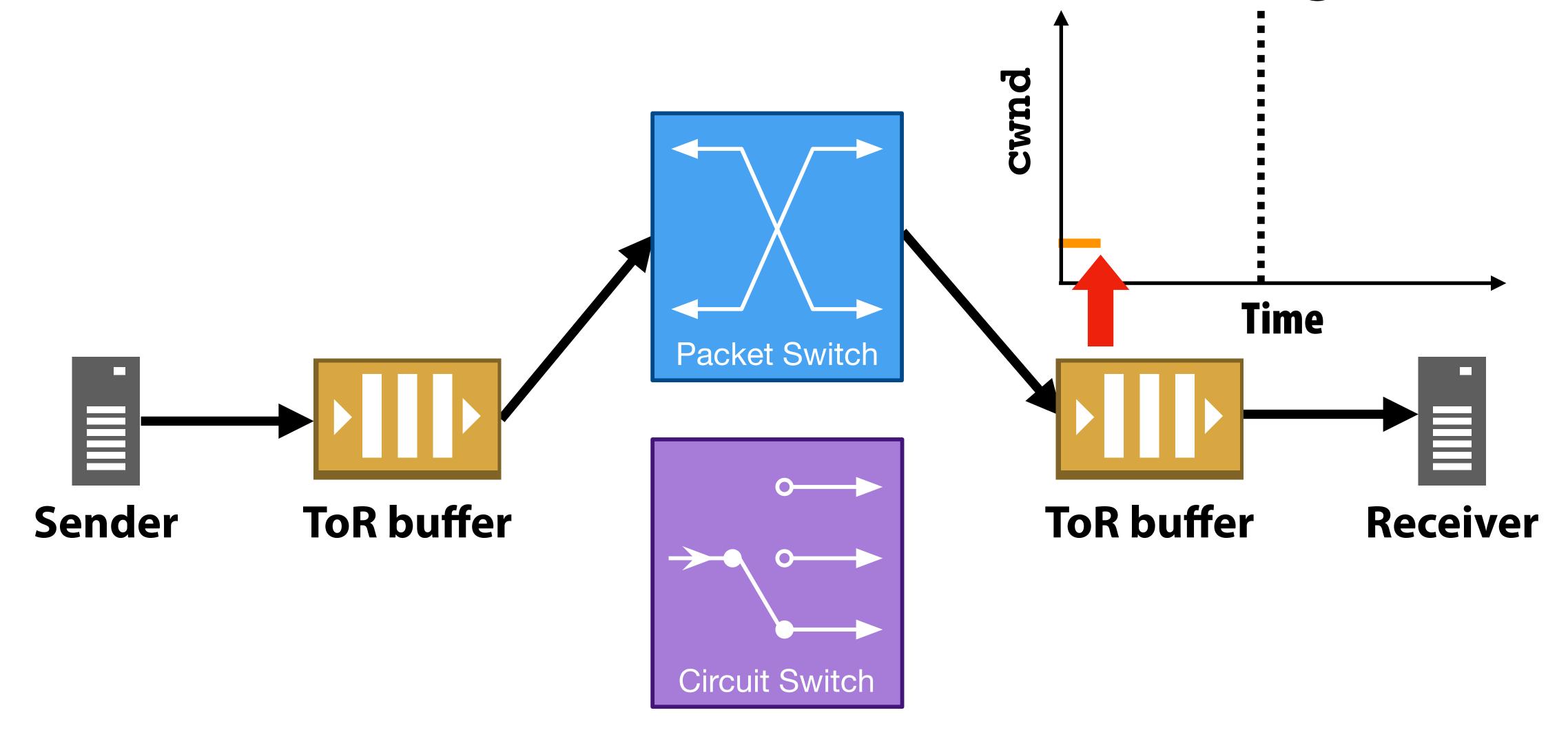
How can we improve this latency?

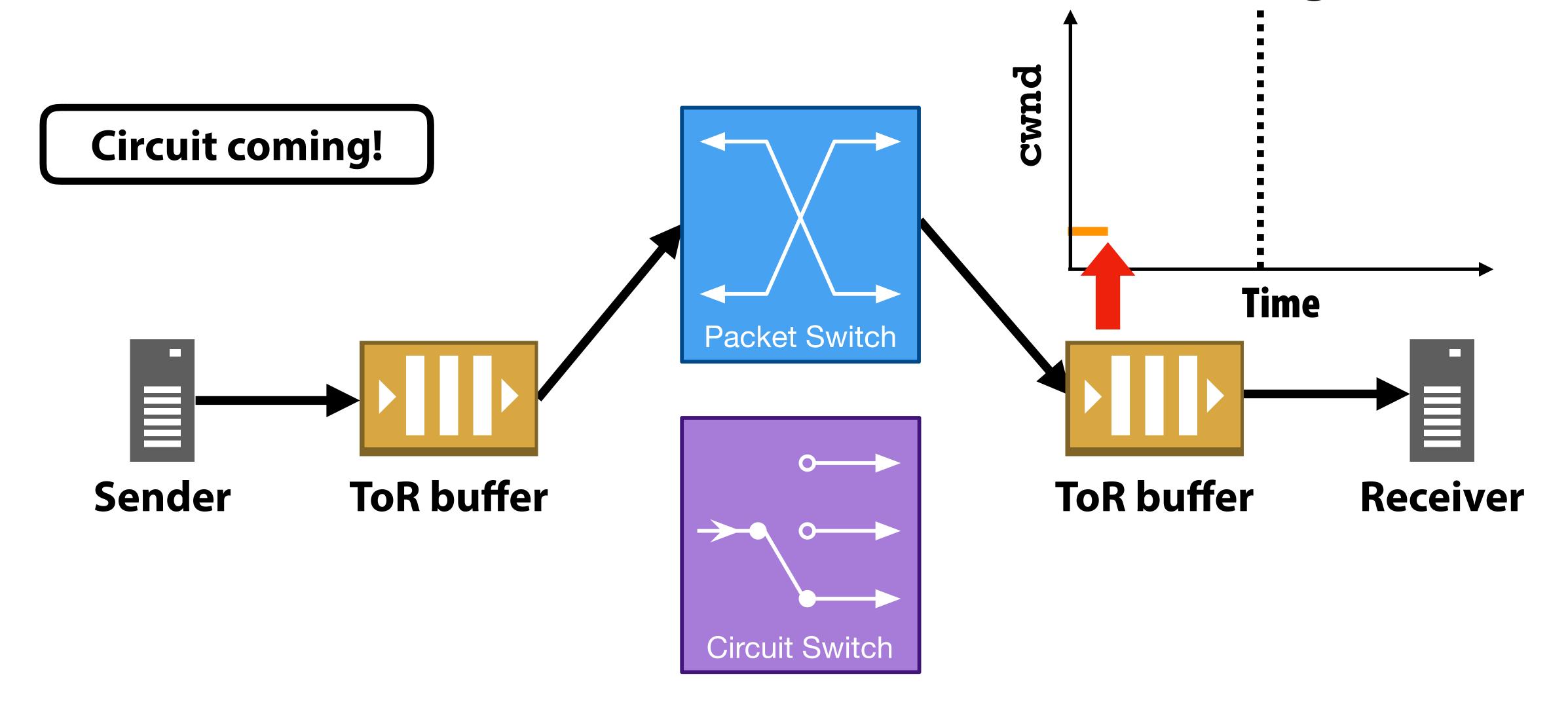
Use large buffers only when circuit is up

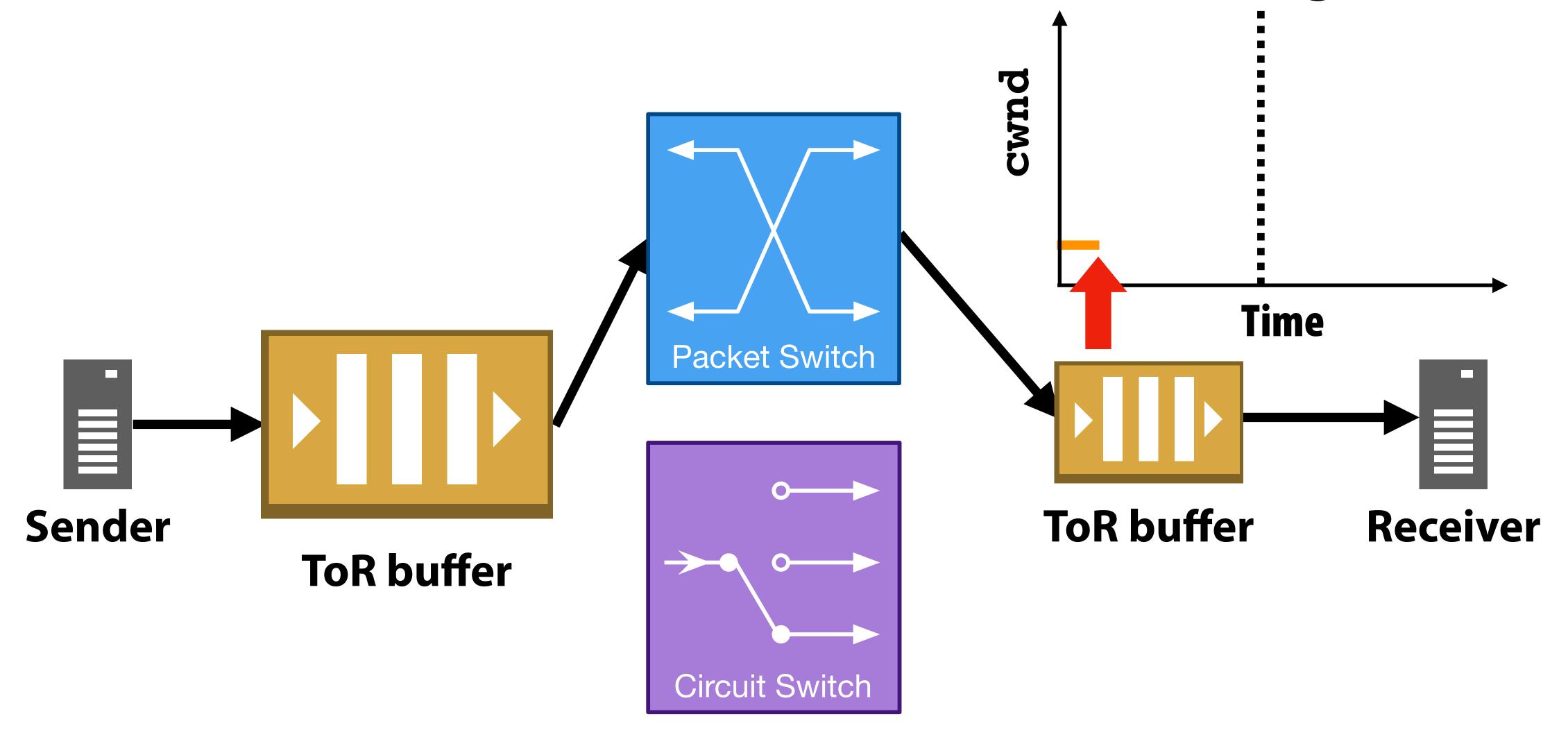
Dynamic buffer resizing: Before a circuit begins, transparently enlarge ToR buffers Full circuit utilization with a latency degradation only during ramp-up period

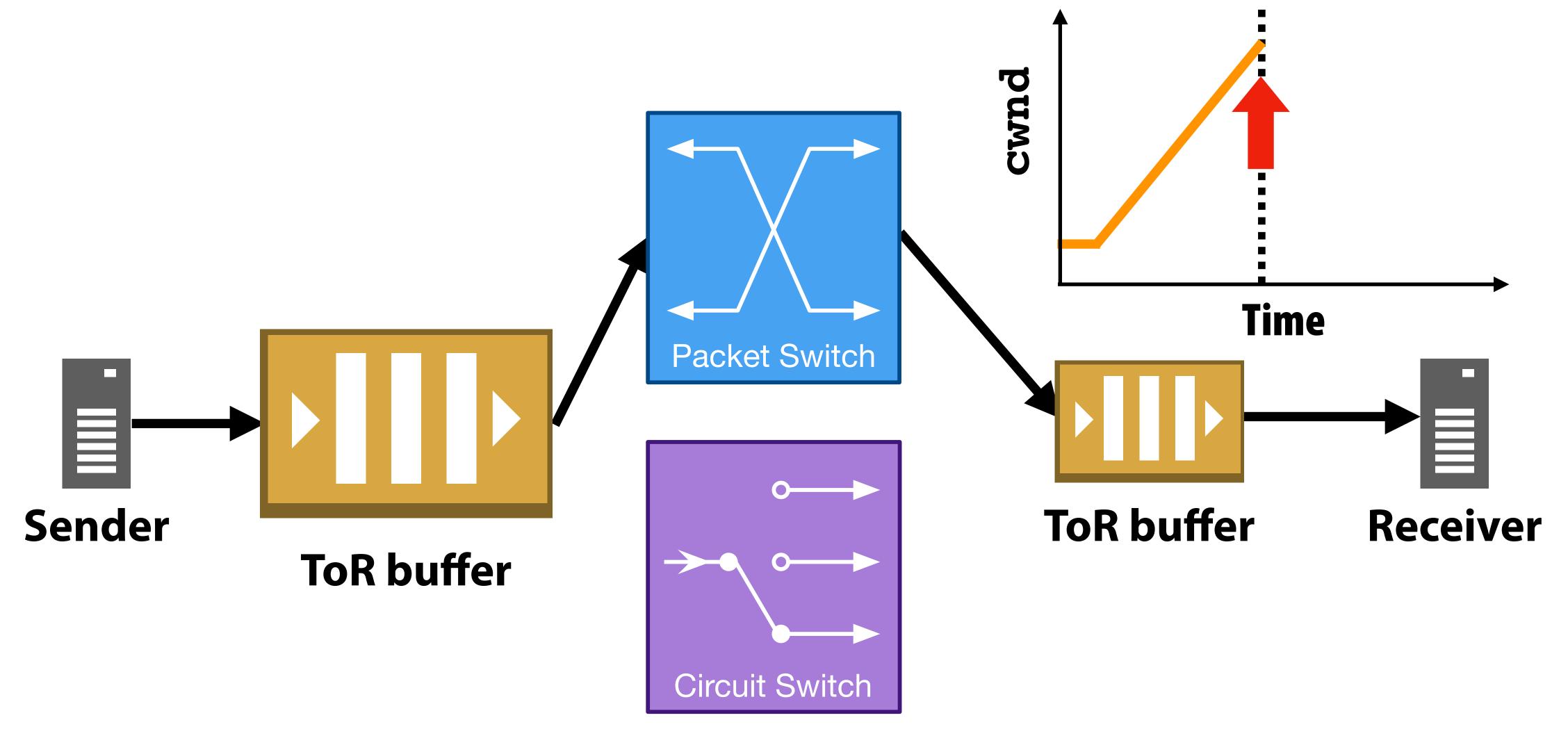


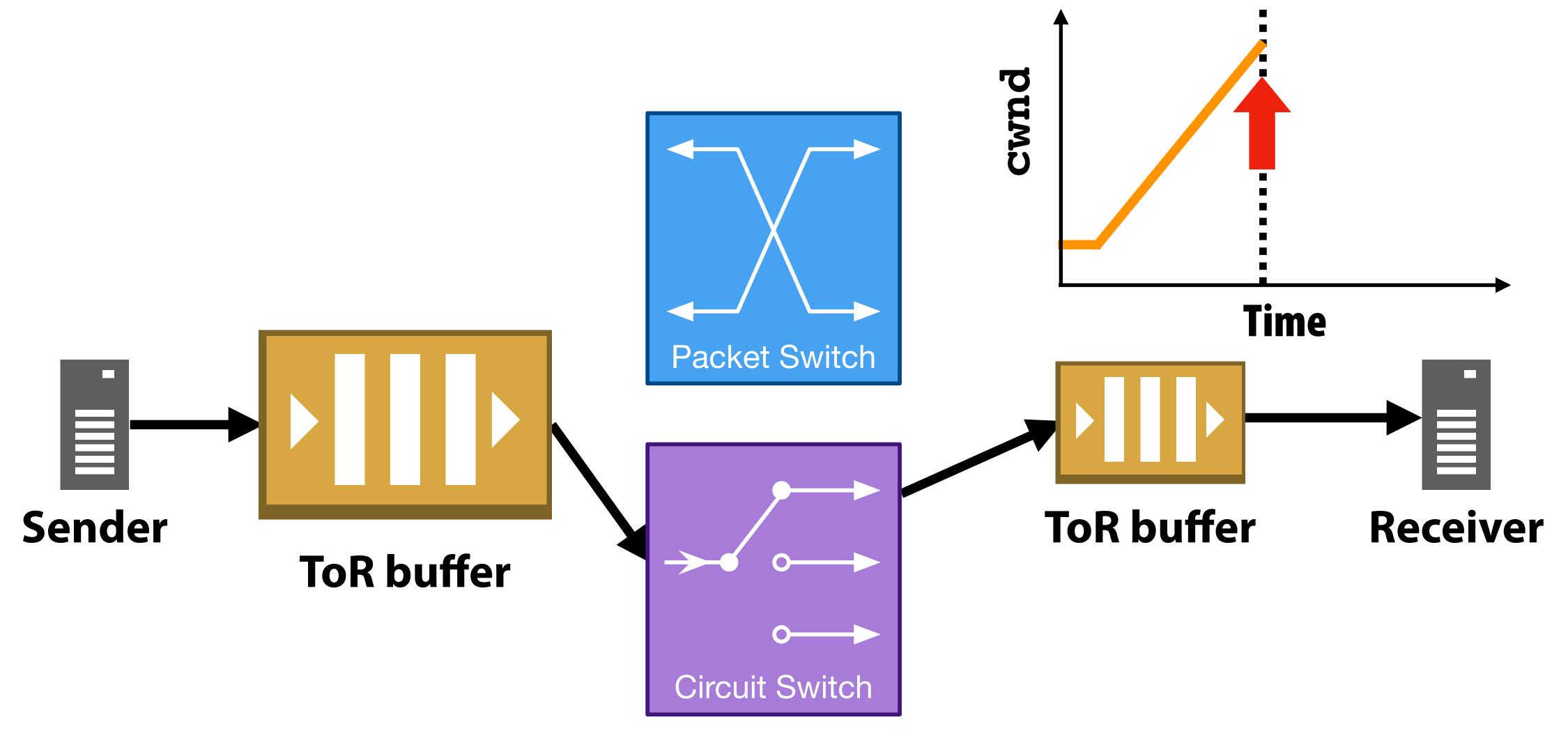


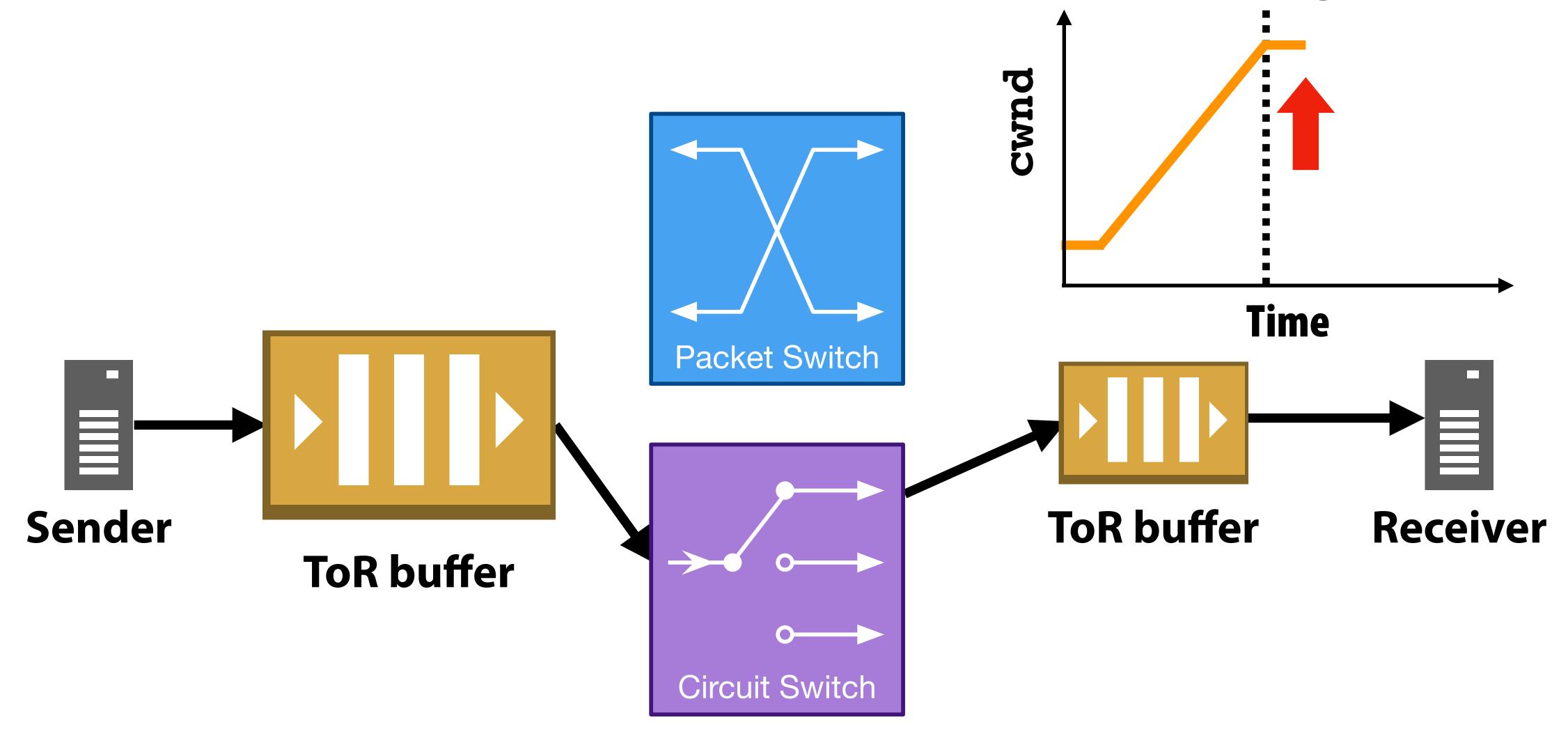


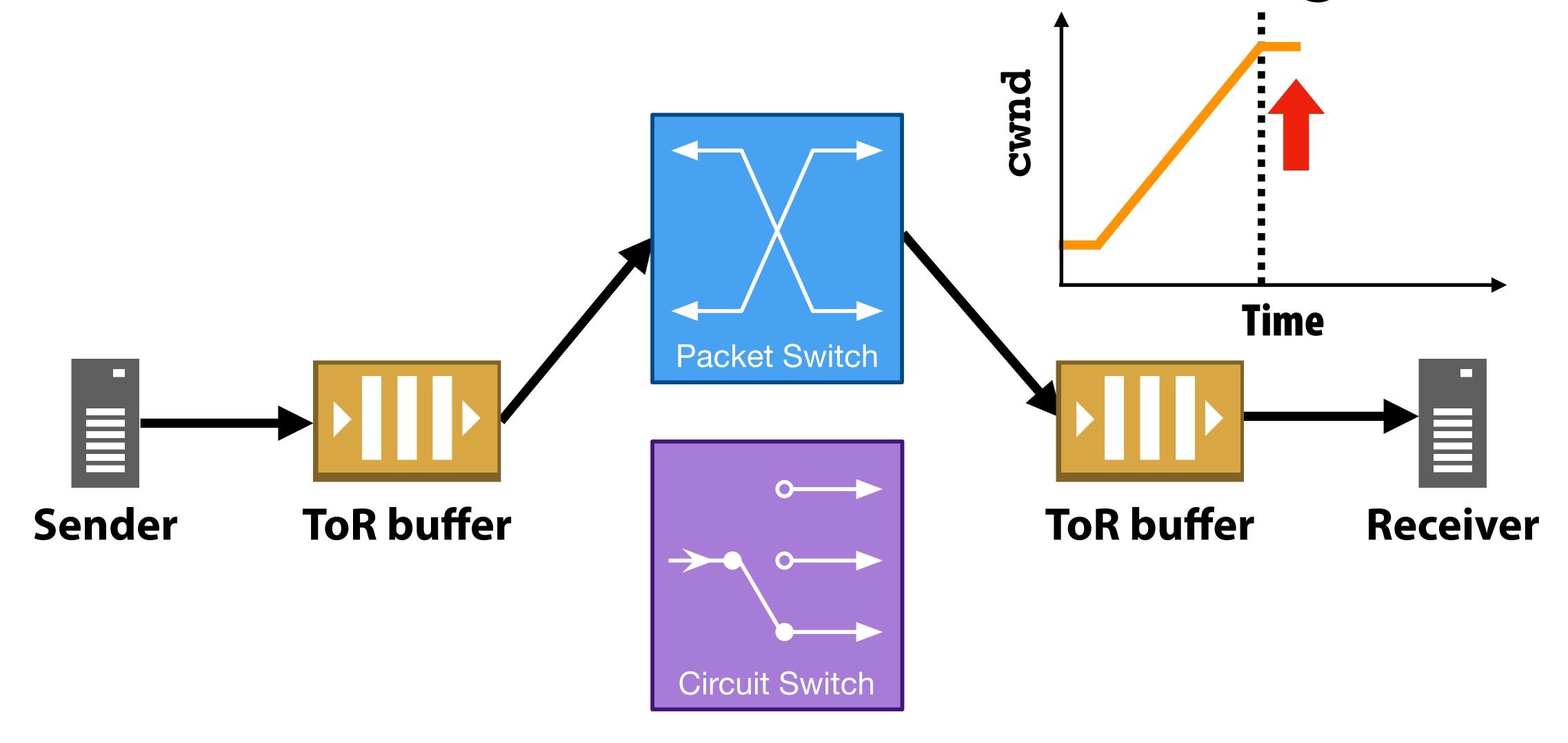


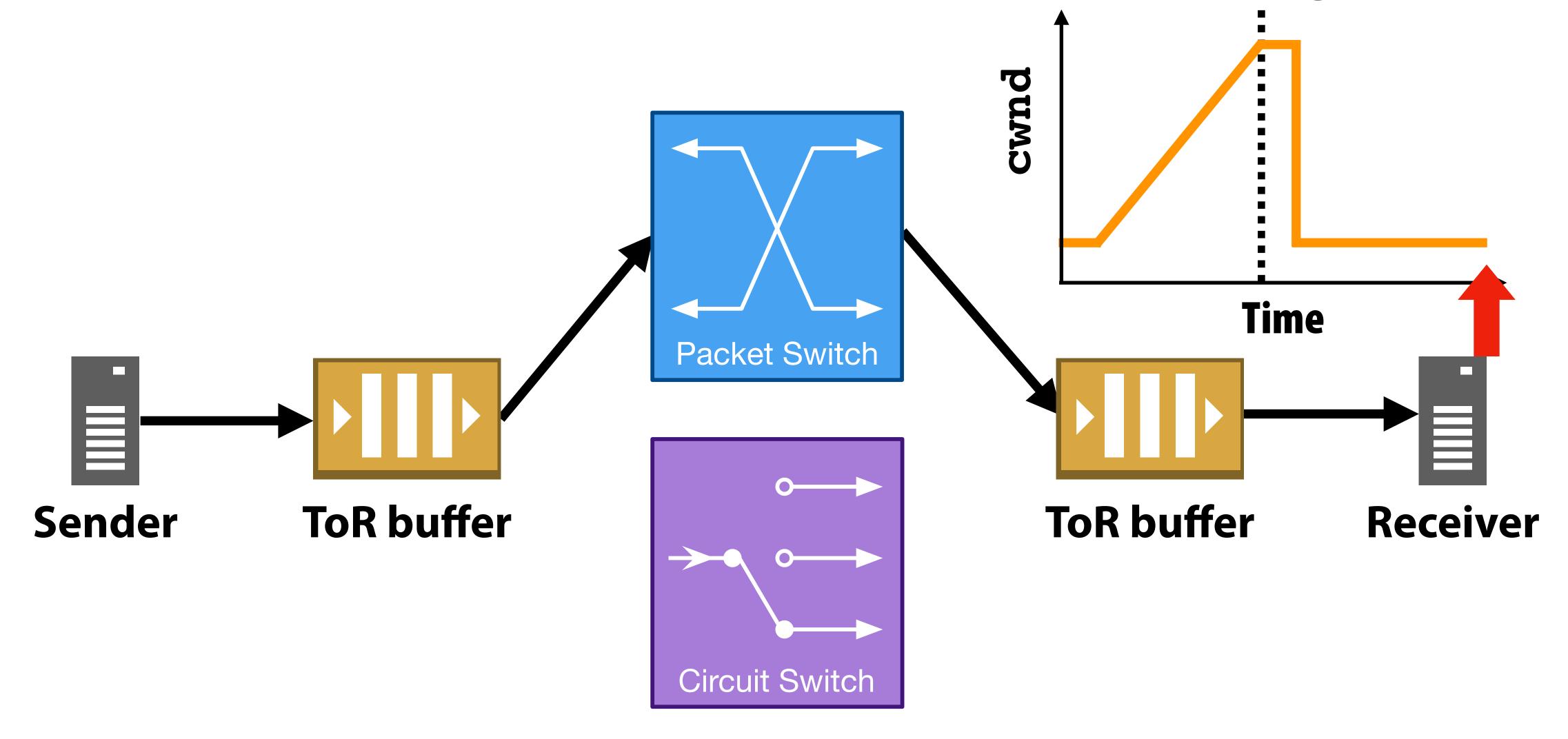












Configuring dynamic buffer resizing

How long in advance should ToR buffers resize (τ) ?

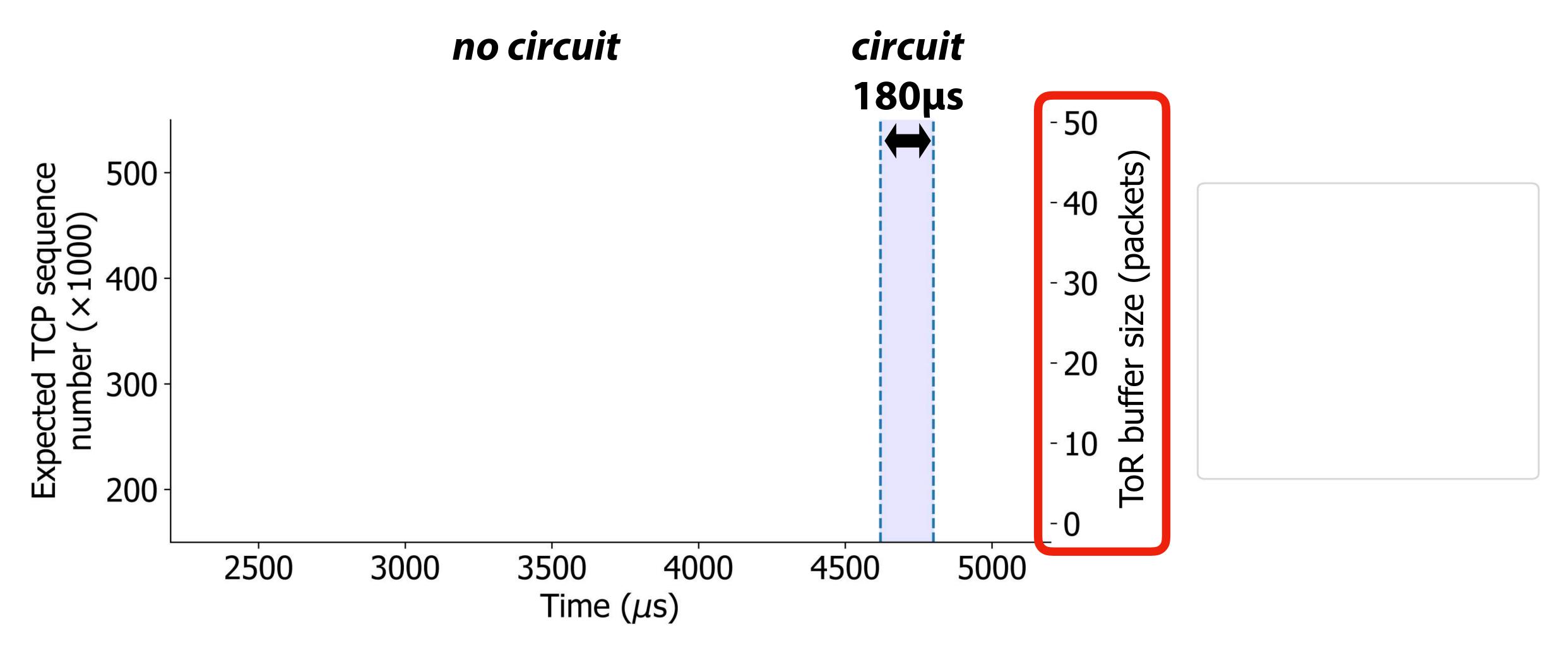
Long enough for TCP to grow cwnd to the circuit BDP

How large should ToR buffers grow to?

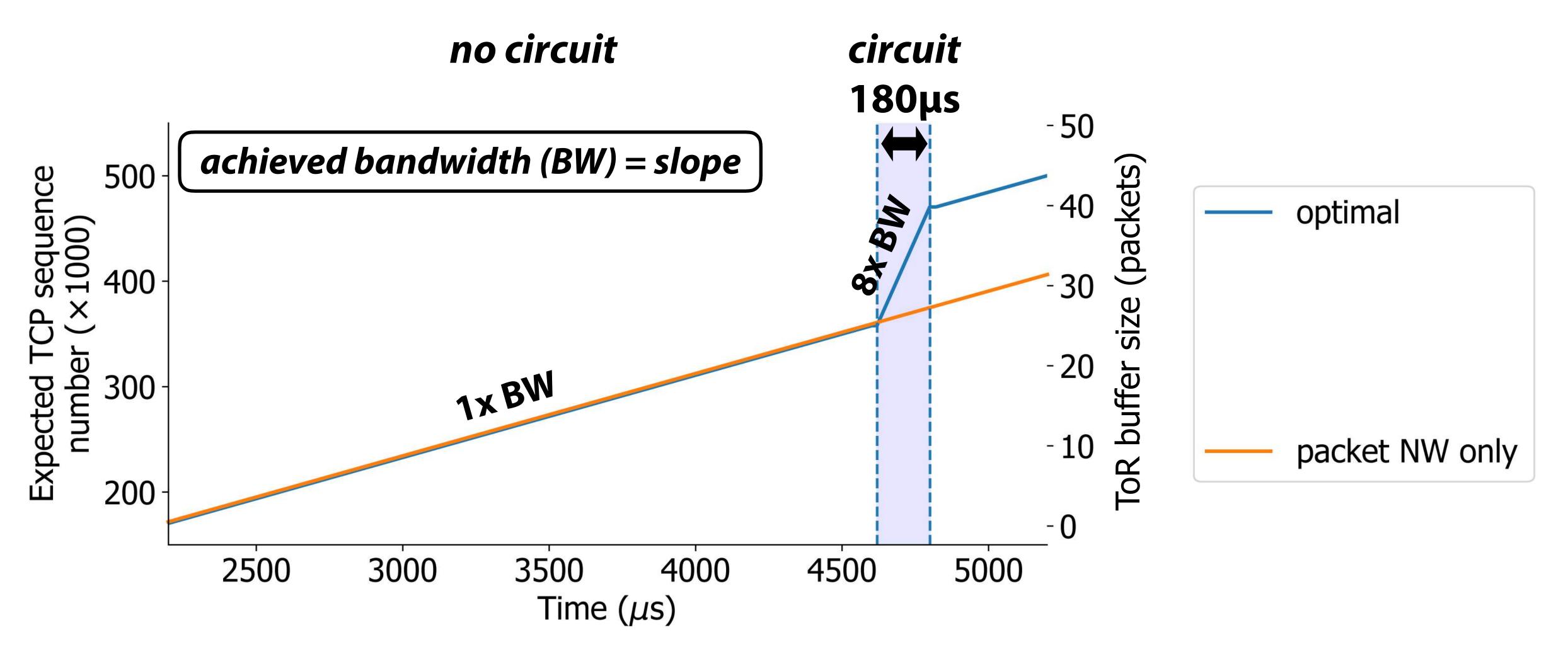
• circuit BDP = 80 Gb/s \times 40 μ s = 45 9000-byte packets

For our configuration, the ToR buffers must hold \sim 40 packets to achieve 90% utilization, which requires 1800 µs of prebuffering

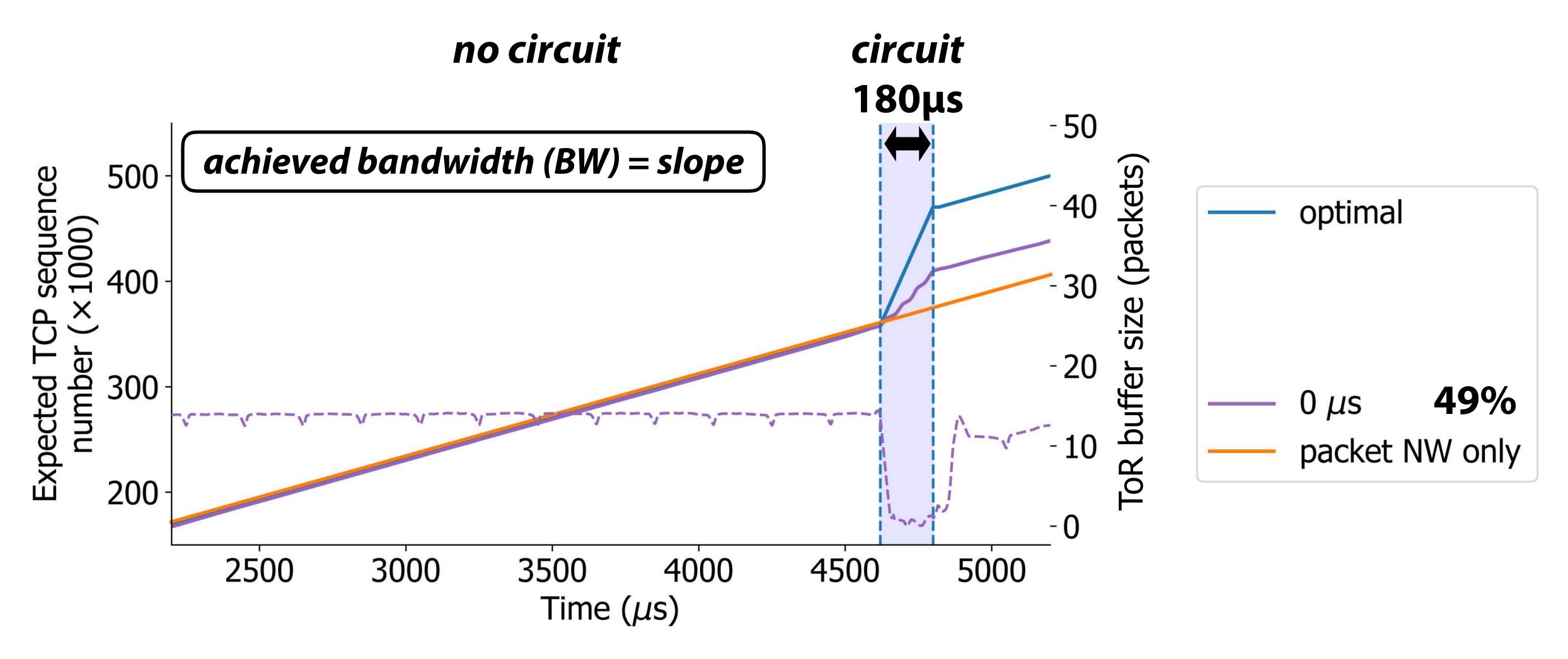
We resize ToR buffers between sizes of 16 and 50 packets



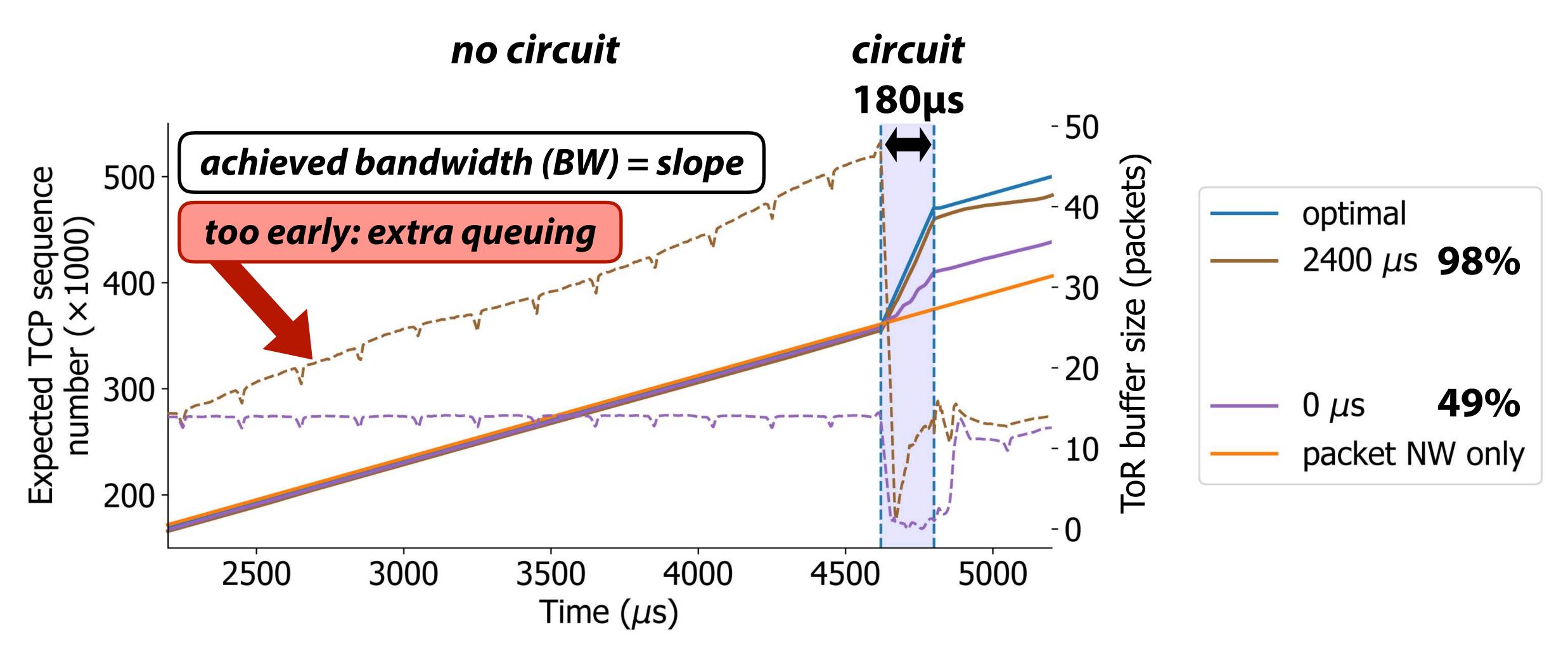
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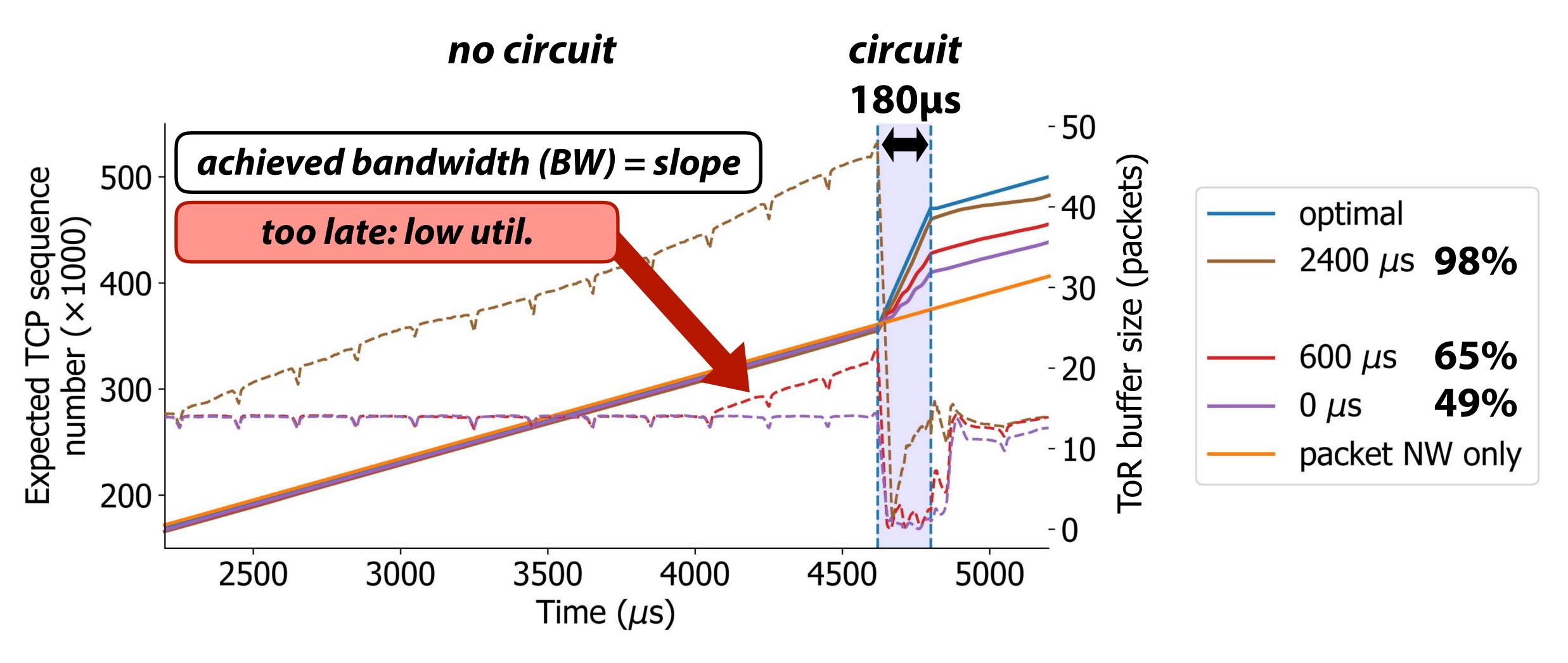


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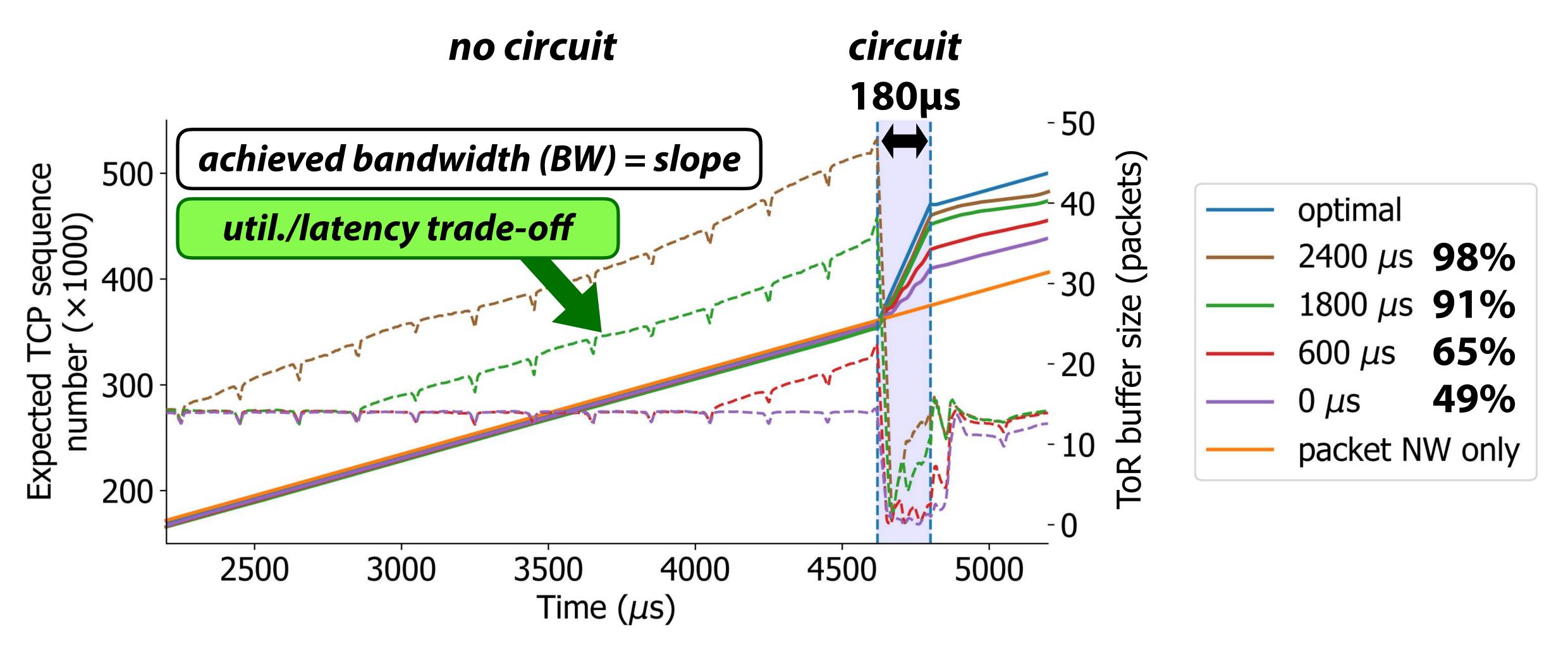
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How long in advance to resize, au?



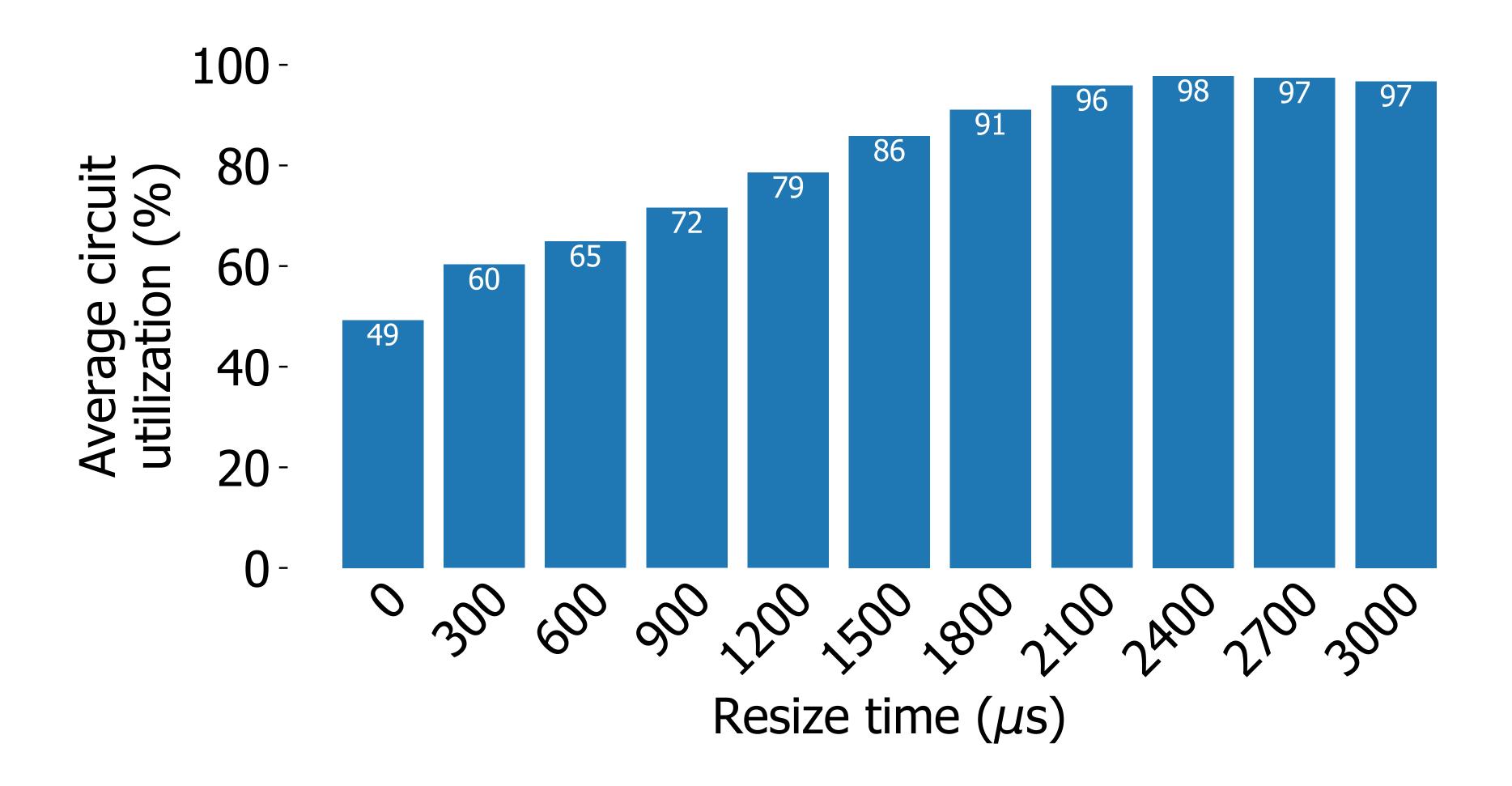
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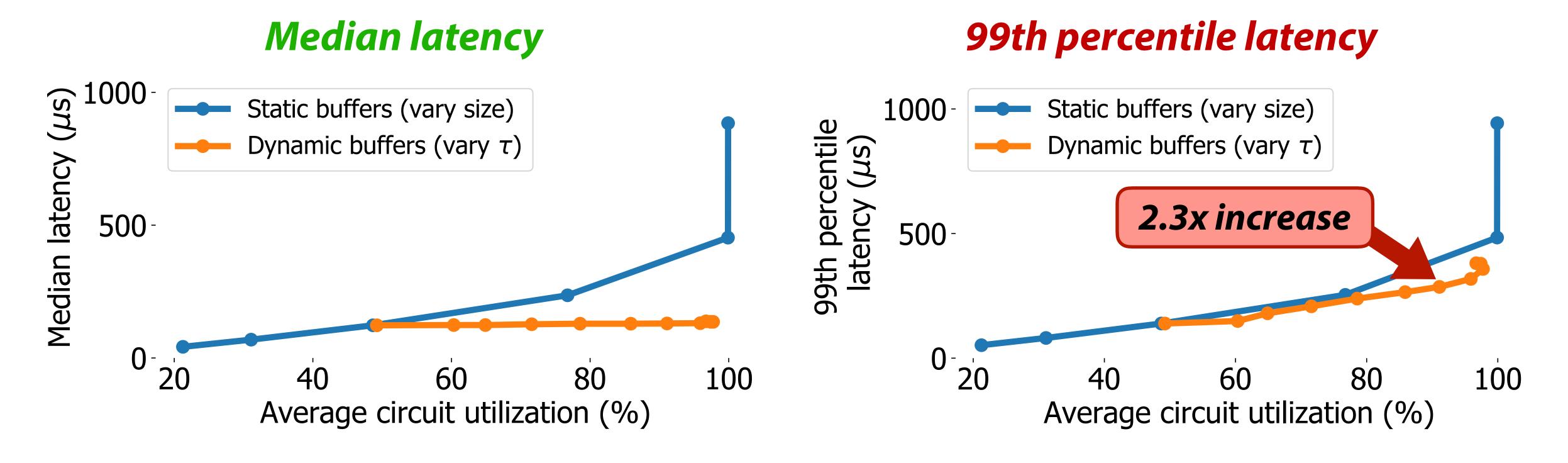
16 flows from rack 1 to rack 2; packet network: 10 Gb/s; circuit network: 80 Gb/s; small buffers: 16 packets; large buffers: 50 packets

1800μs of prebuffering yields 91% util.



16 flows from rack 1 to rack 2; packet network: 10 Gb/s; circuit network: 80 Gb/s; small buffers: 16 packets; large buffers: 50 packets

Latency degradation during ramp-up



We cannot use large queues for so long.

Can we get the same high utilization with shorter prebuffering?

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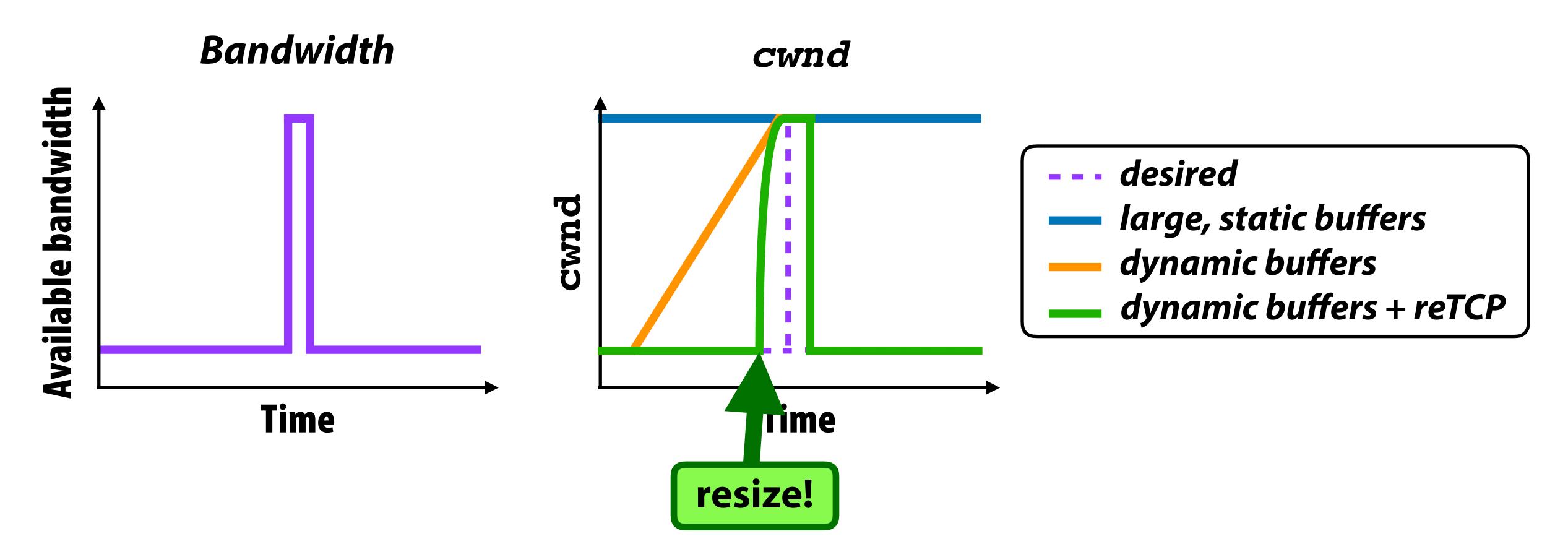
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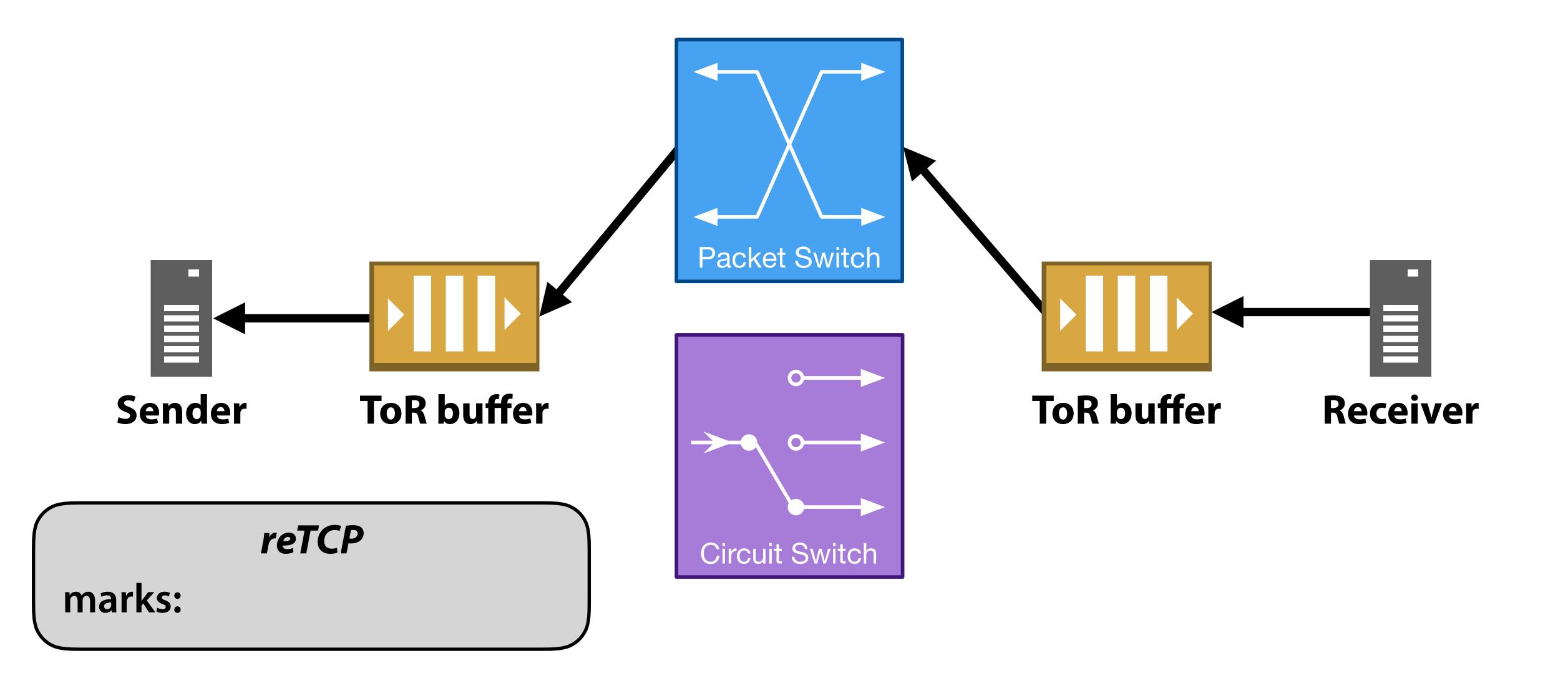
Mitigates tail latency penalty

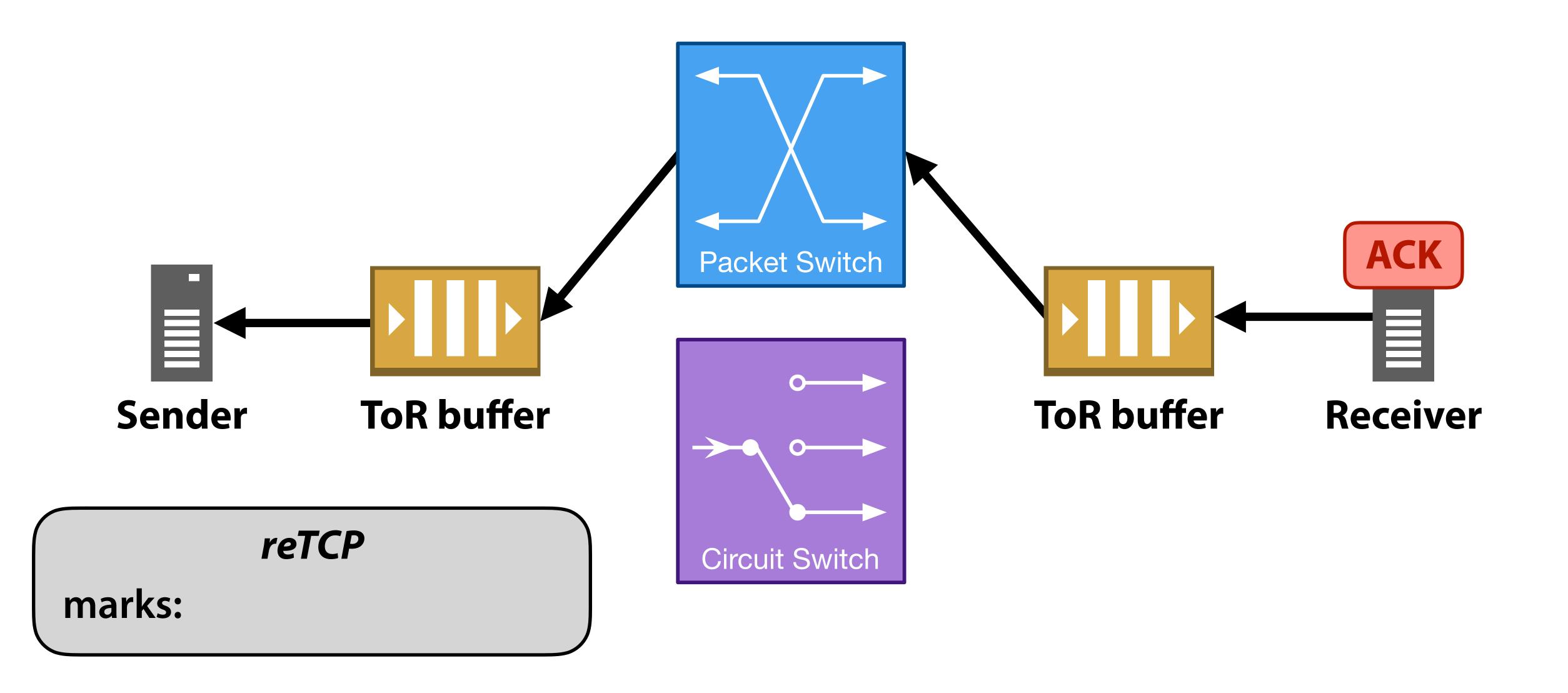
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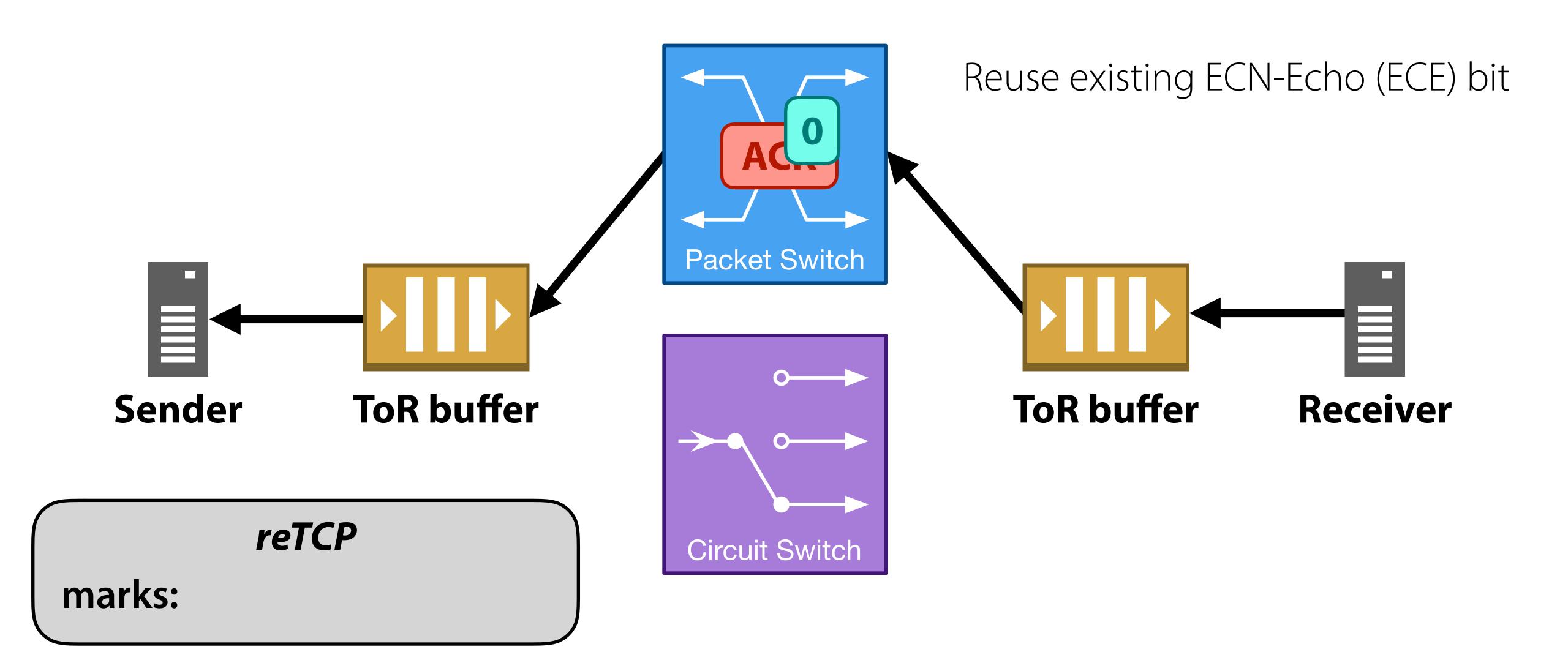
reTCP: Rapidly grow cwnd before a circuit

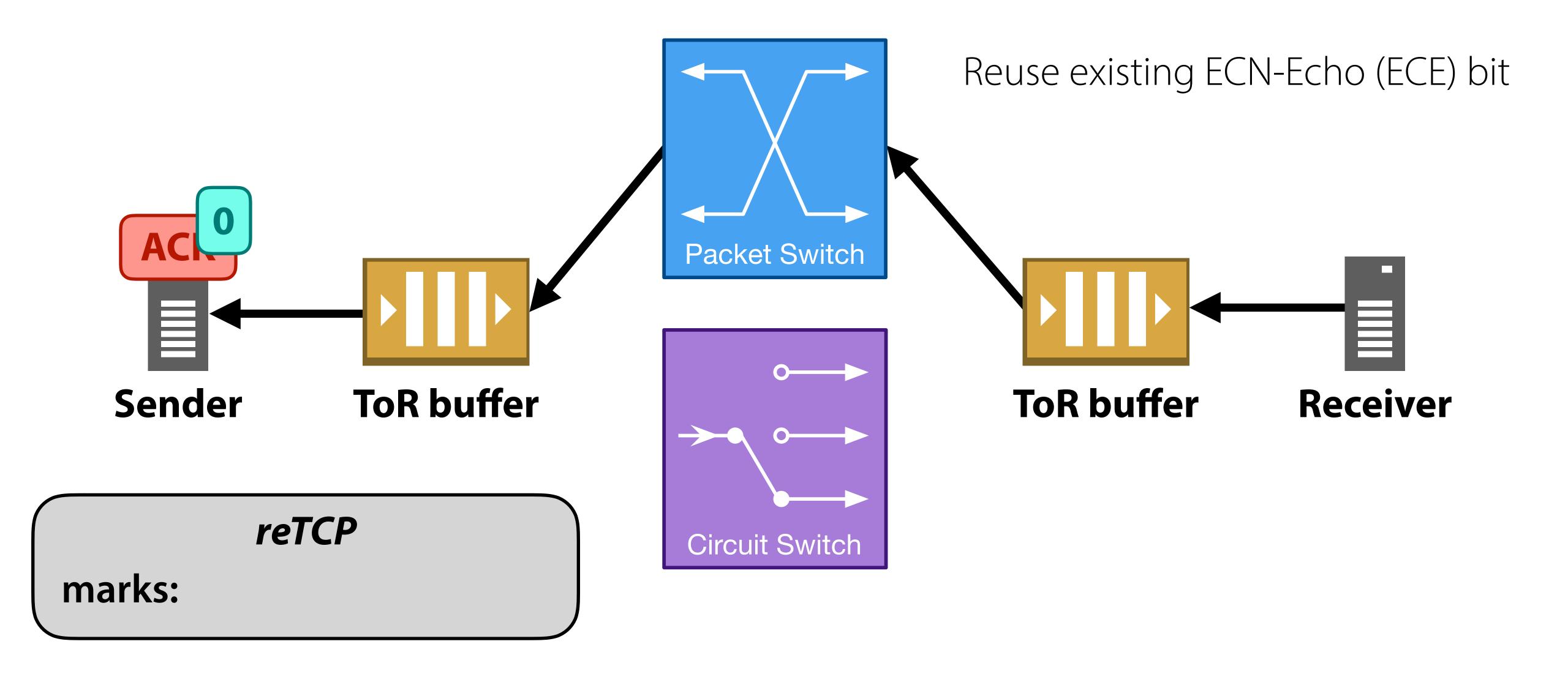
- 1) Communicate circuit state to sender TCP
- 2) Sender TCP reacts by multiplicatively increasing/decreasing cwnd

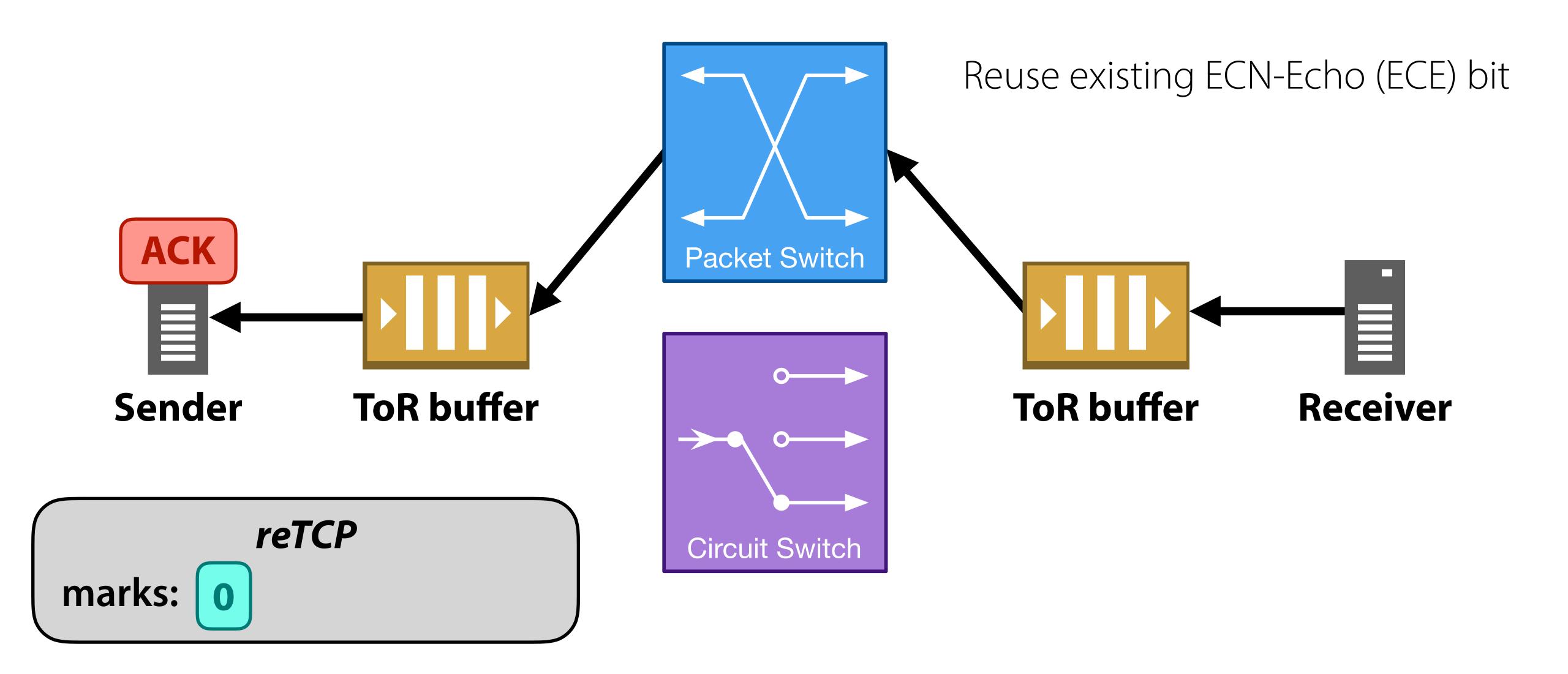


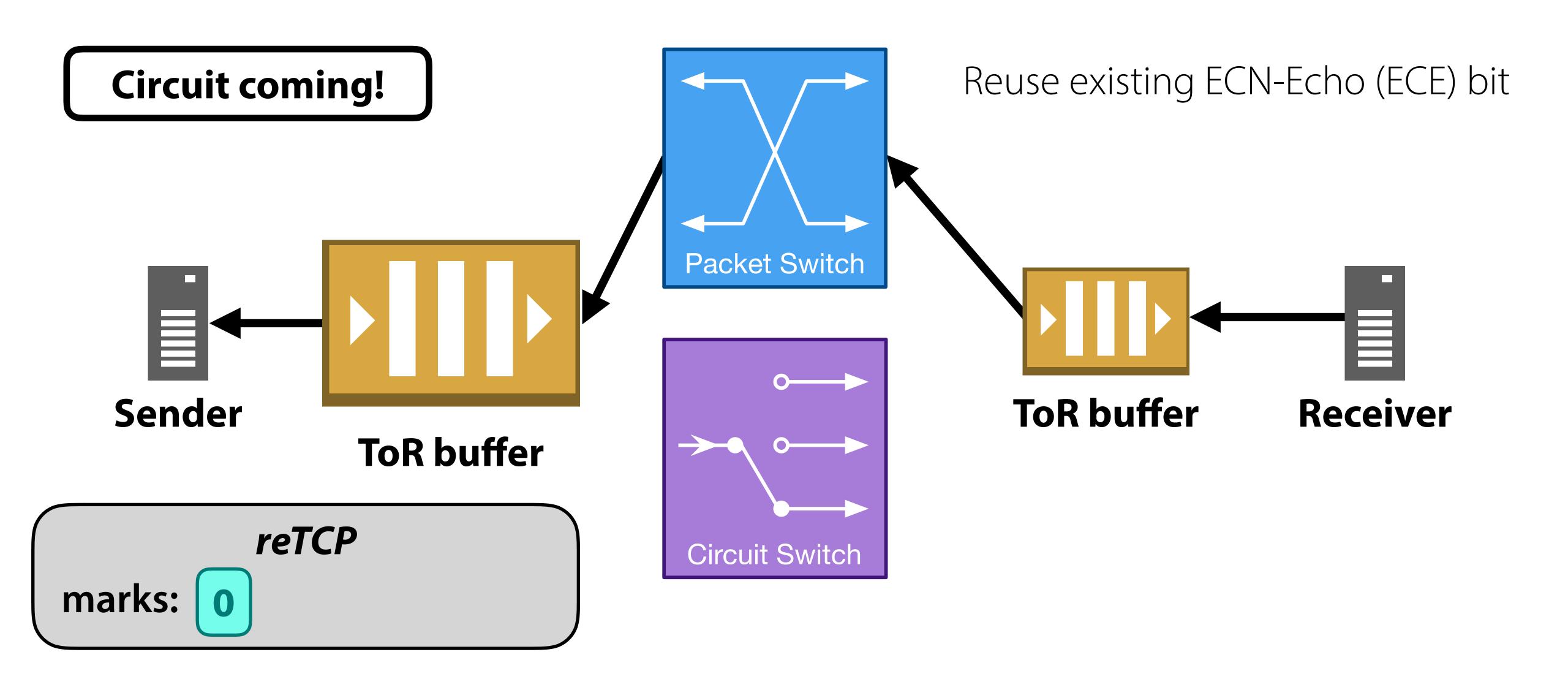


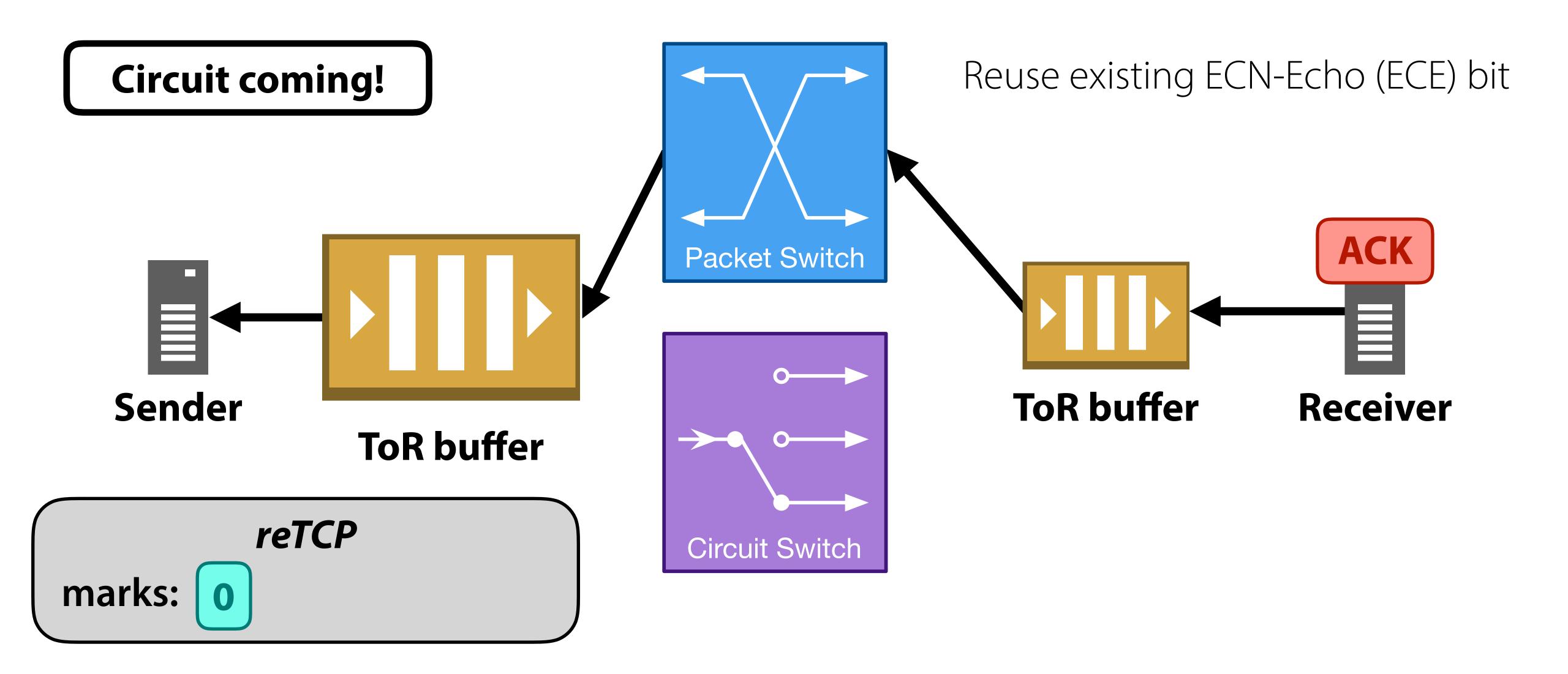


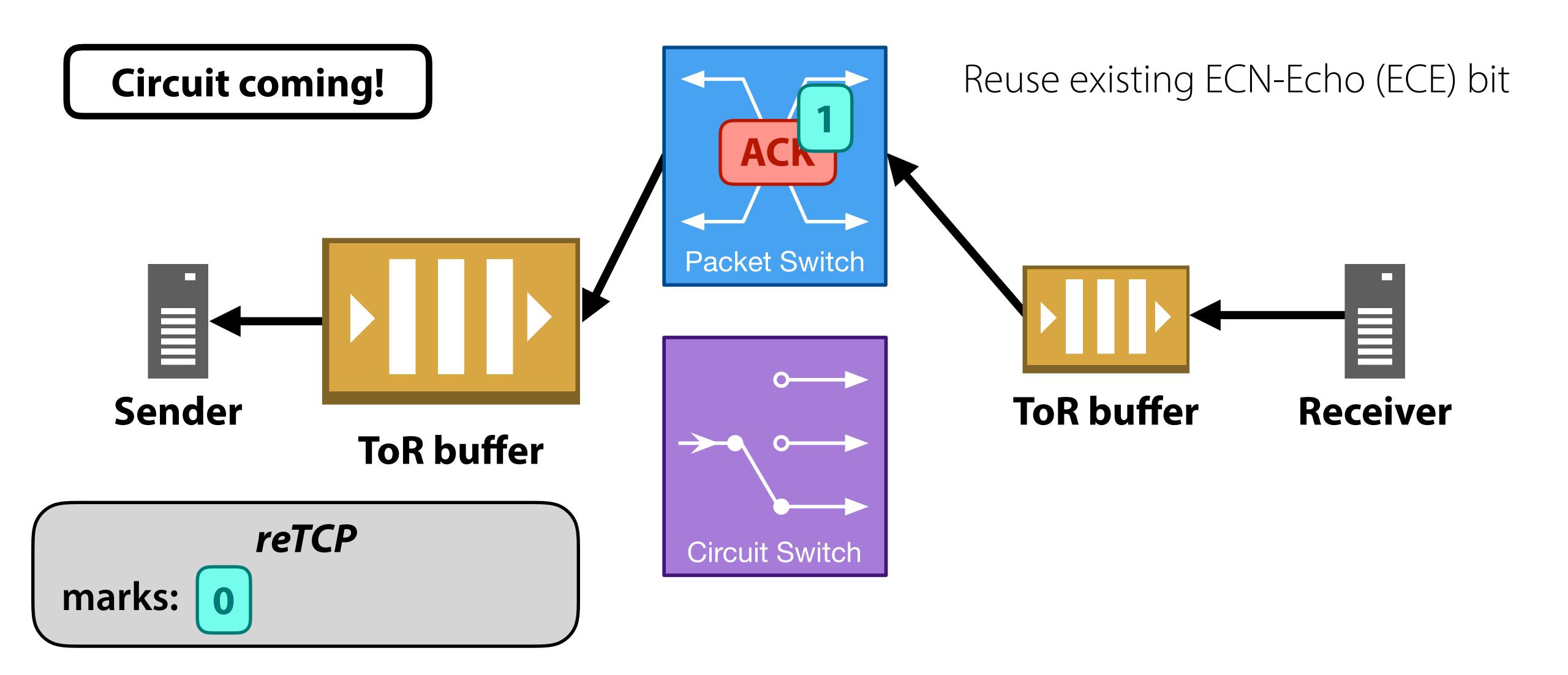


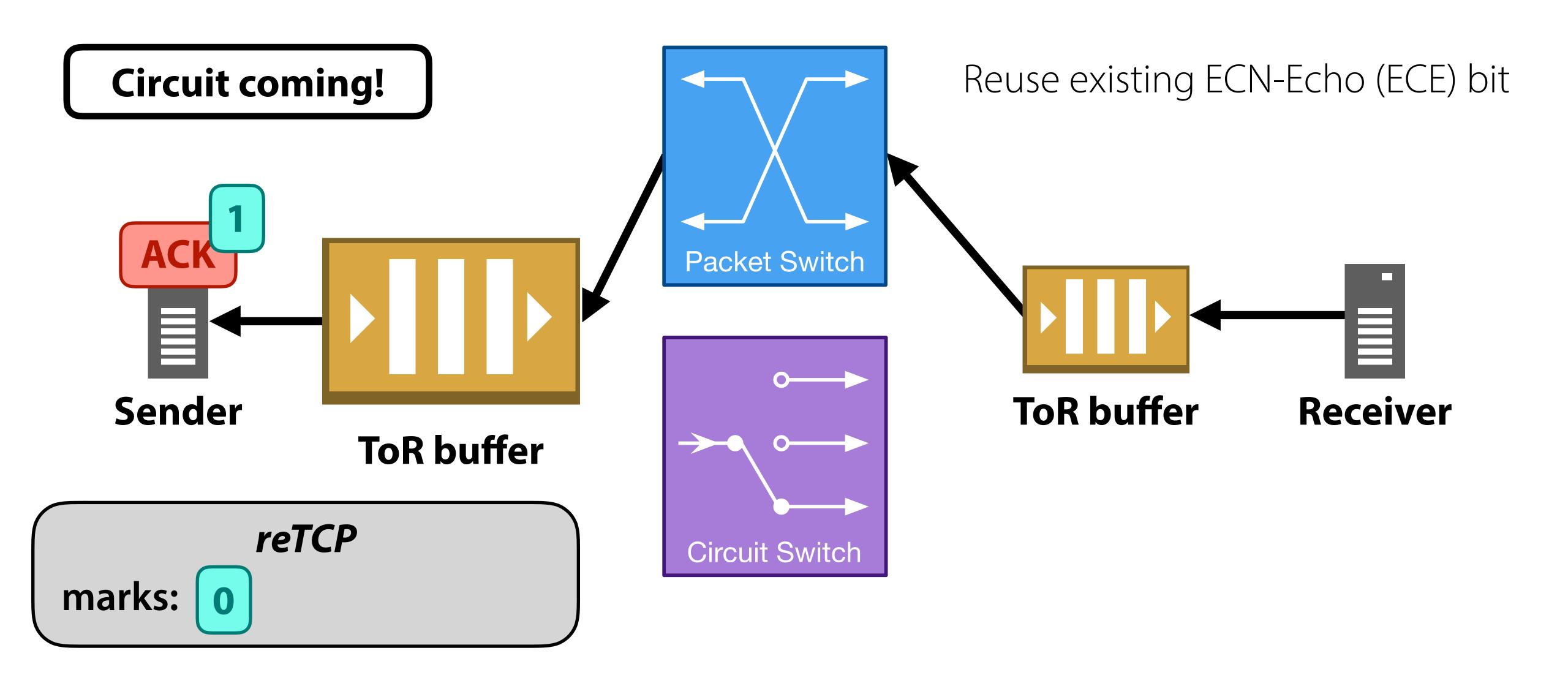


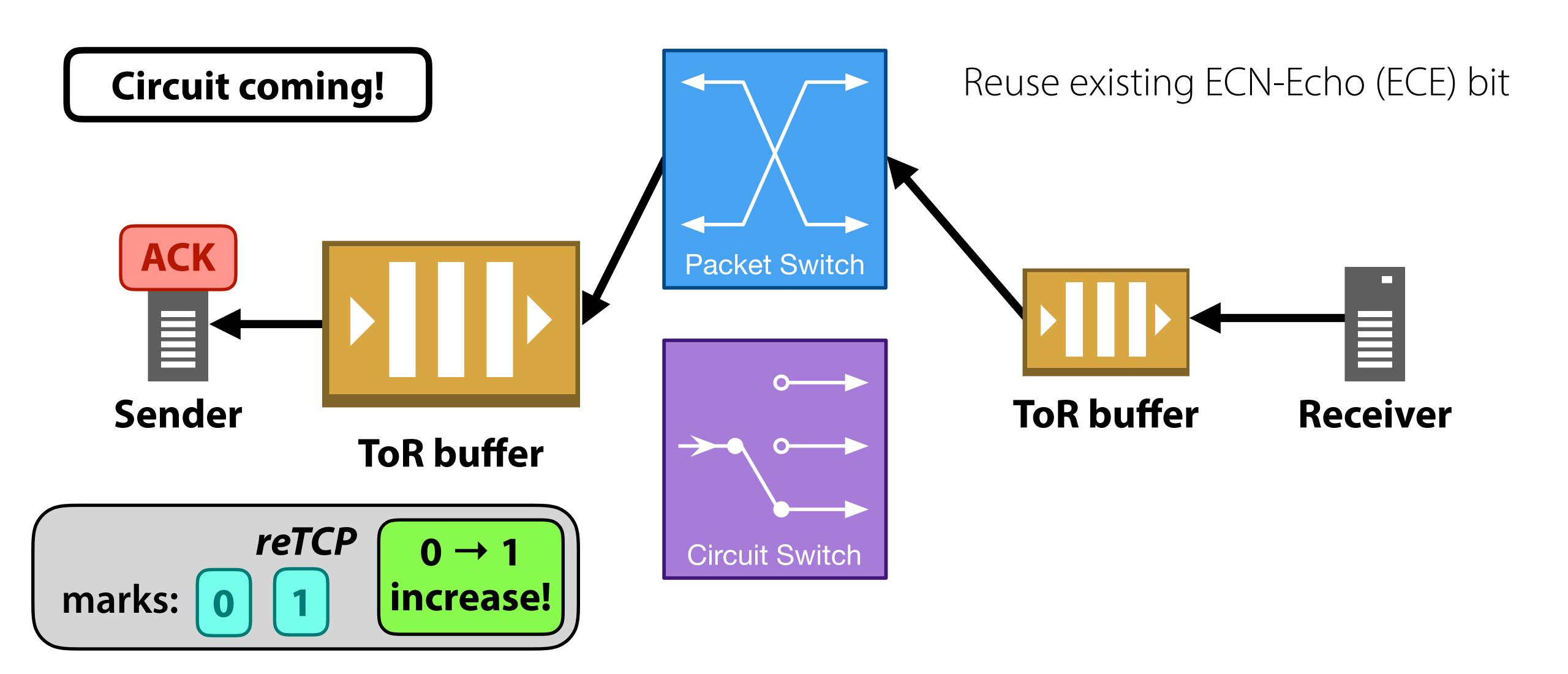


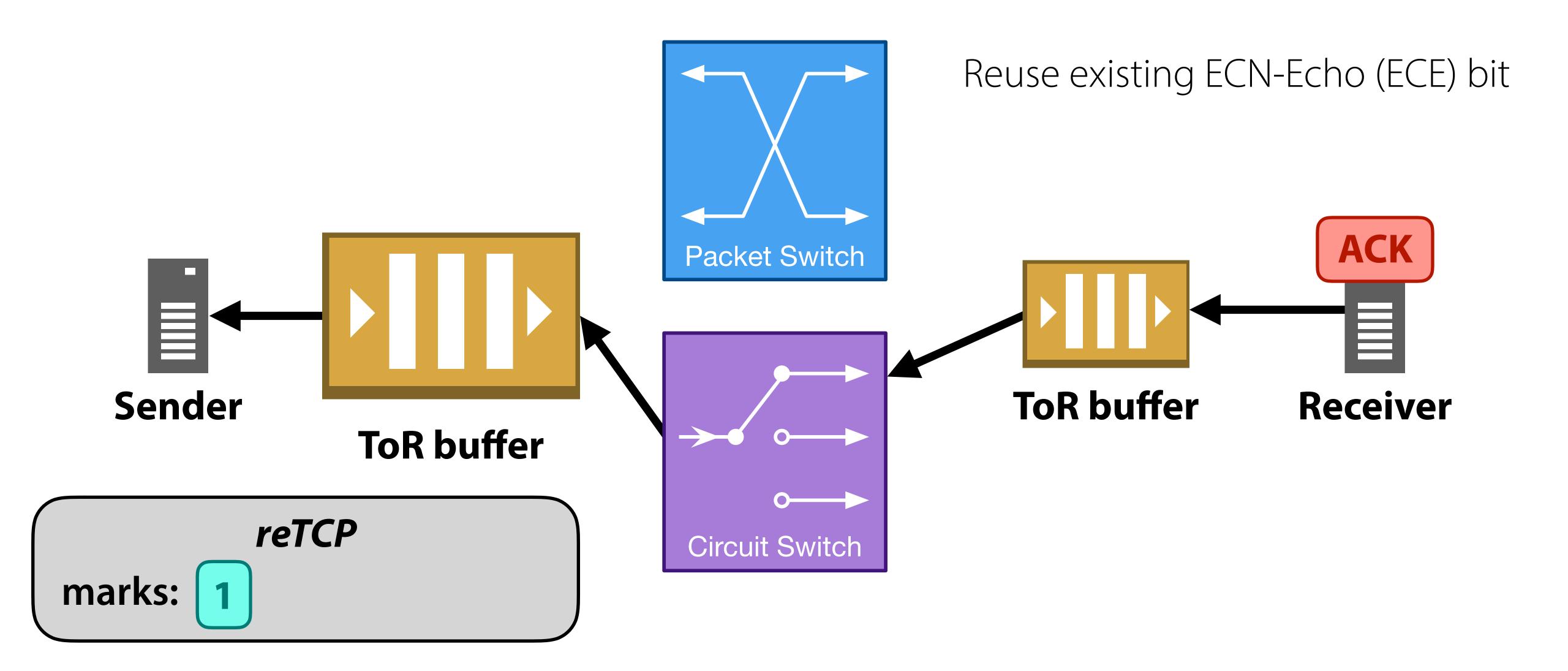


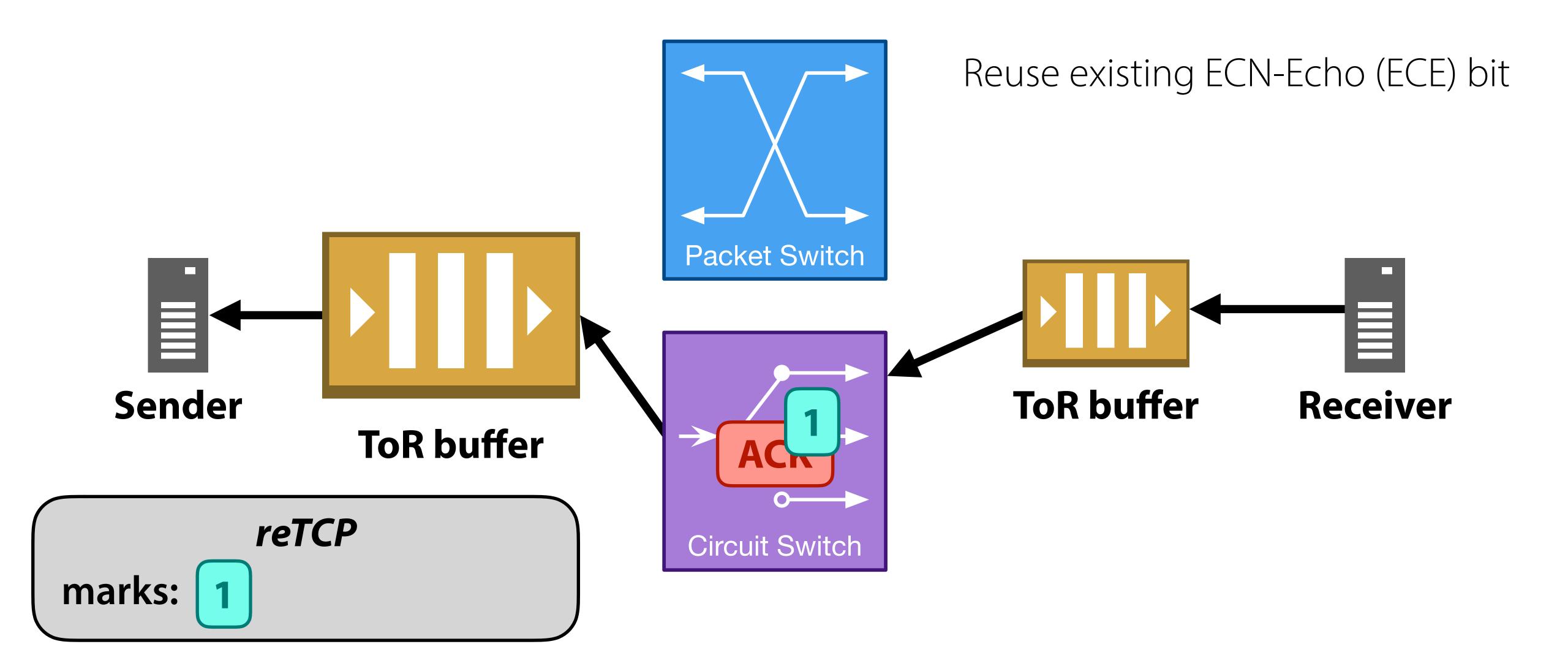


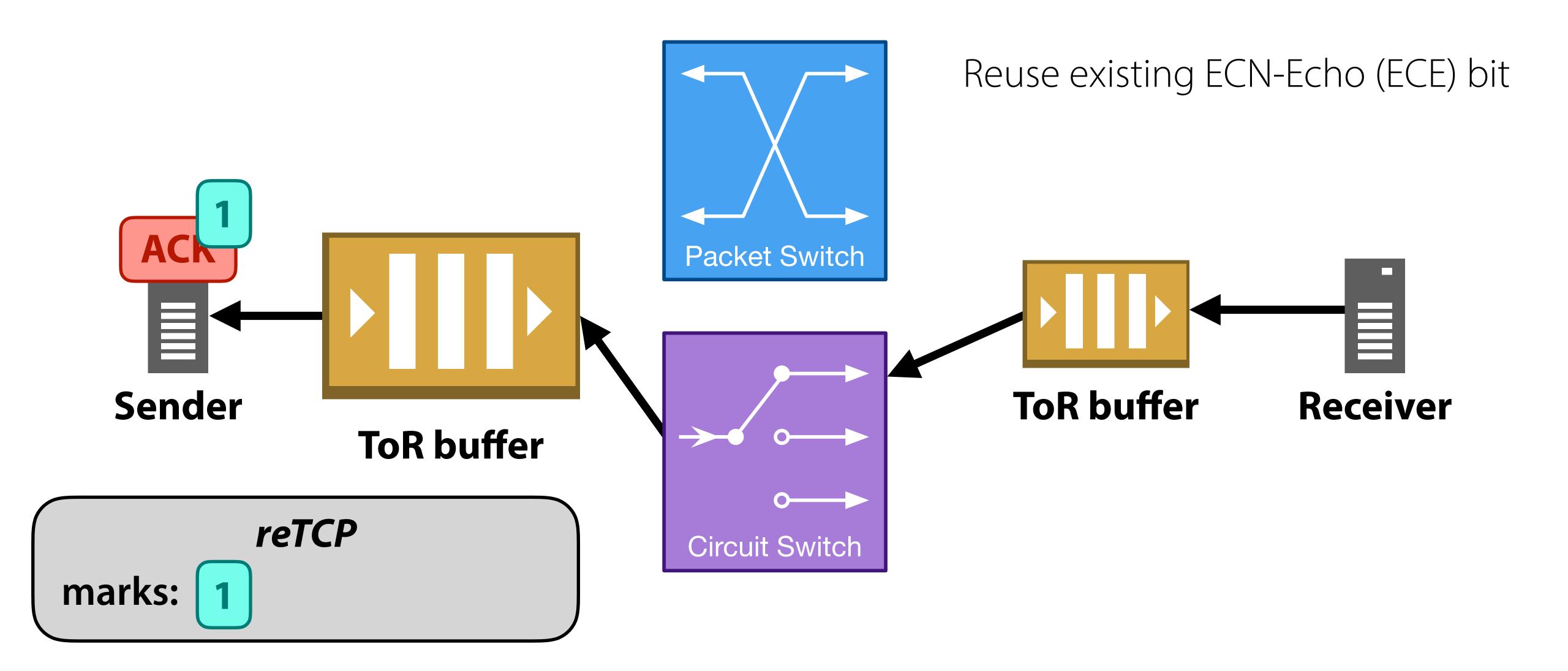


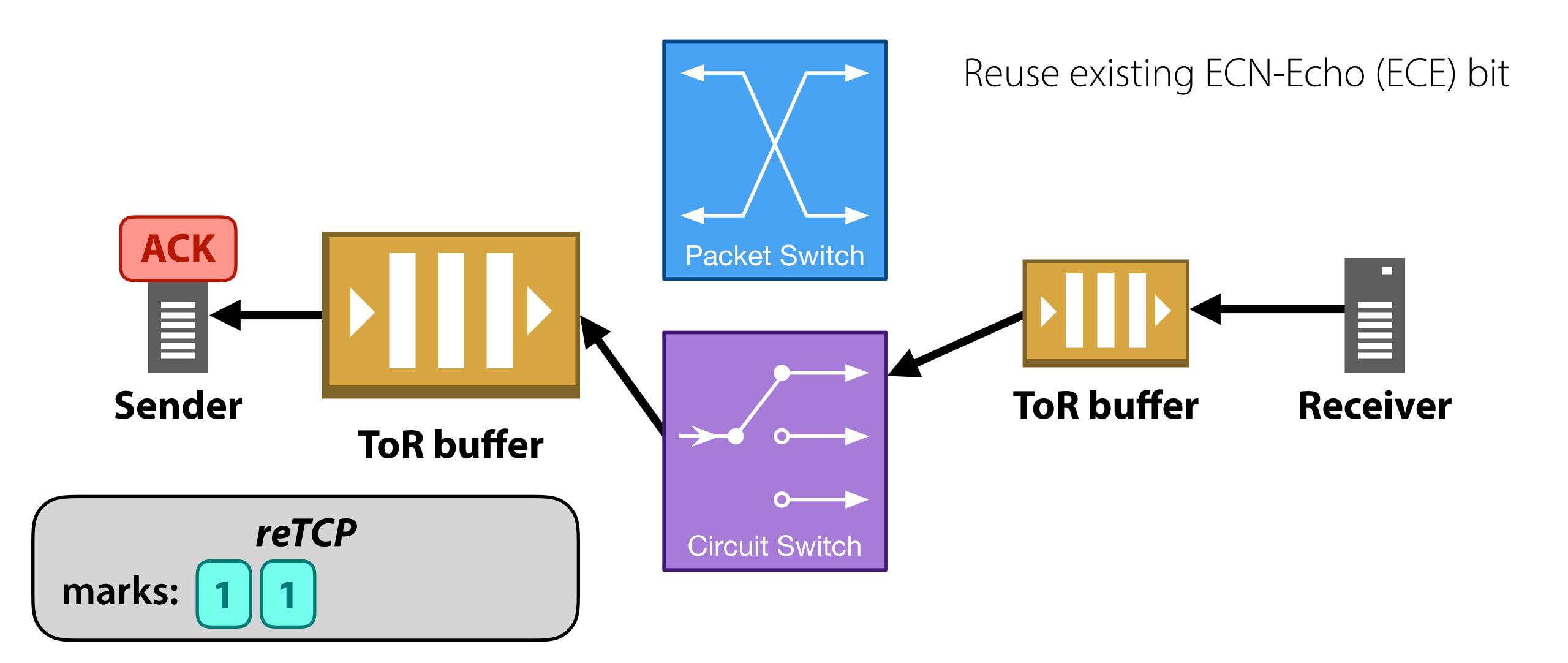


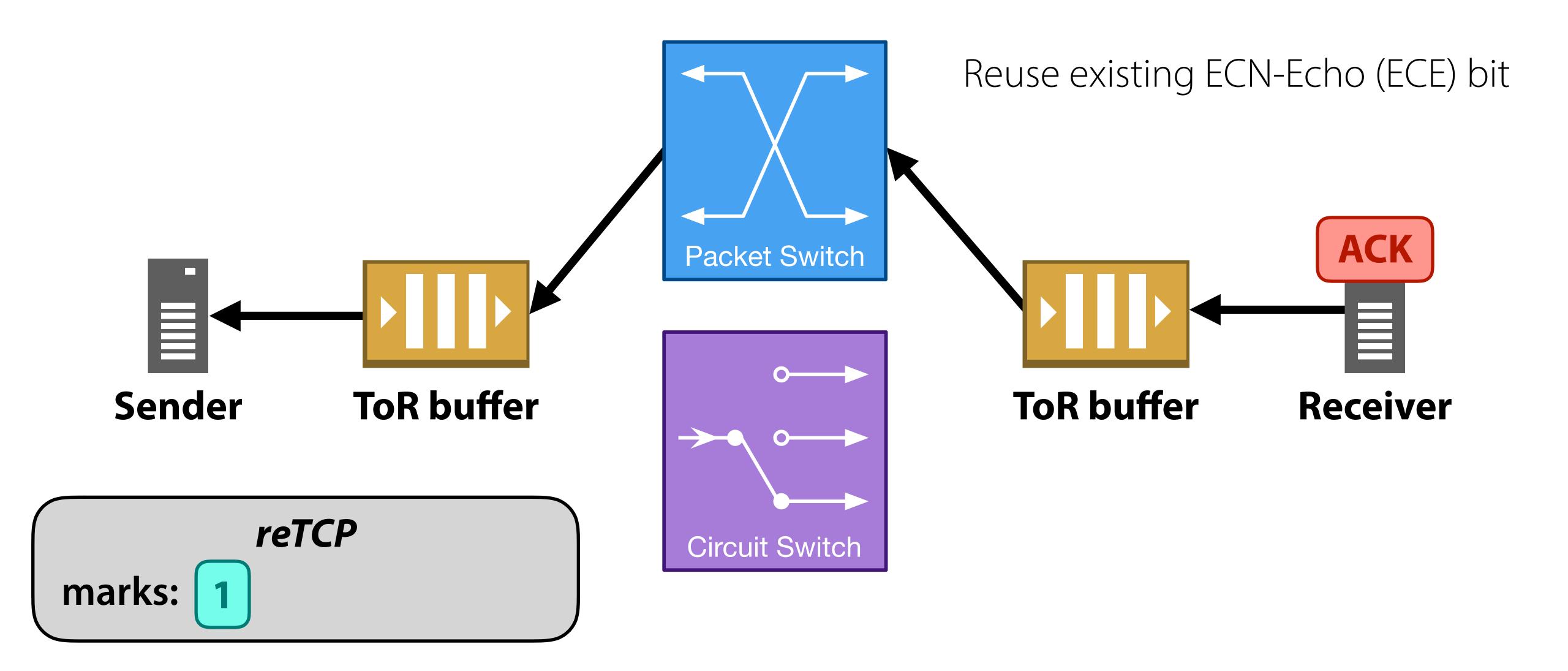


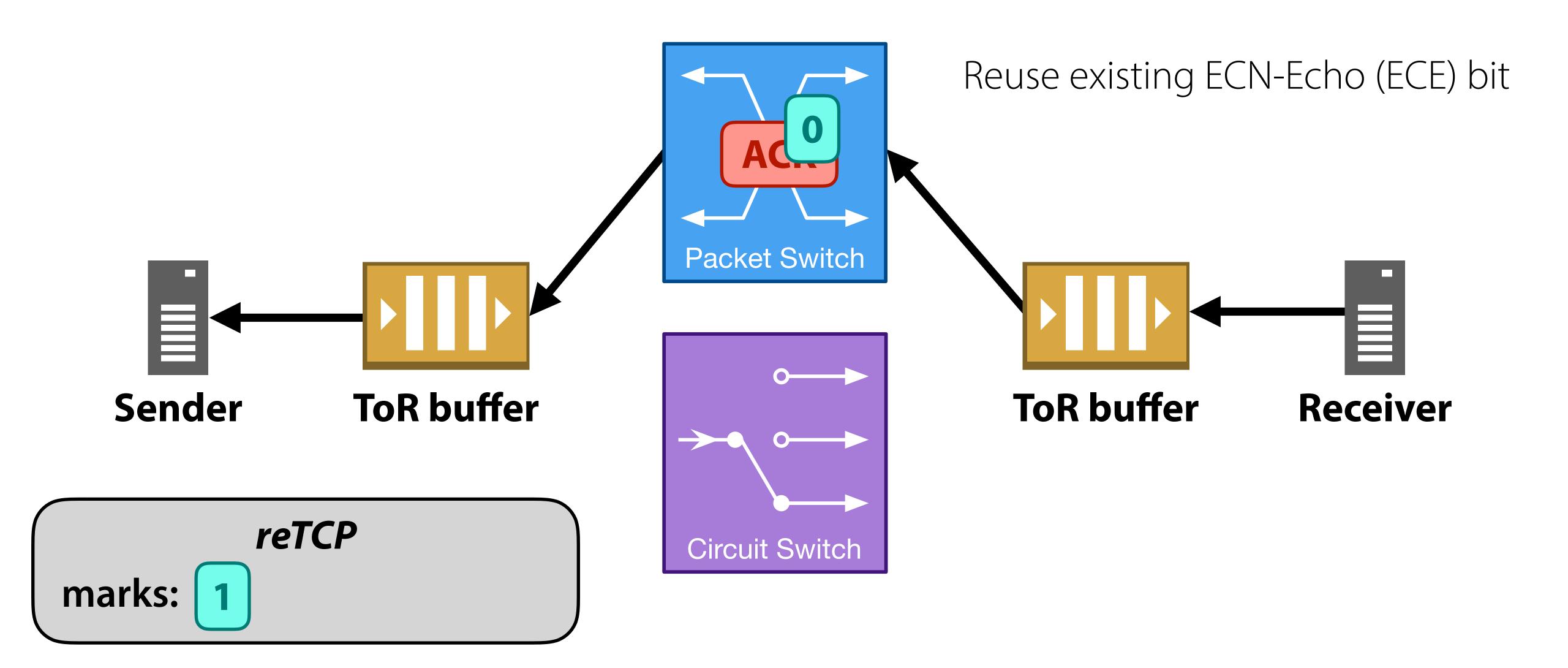


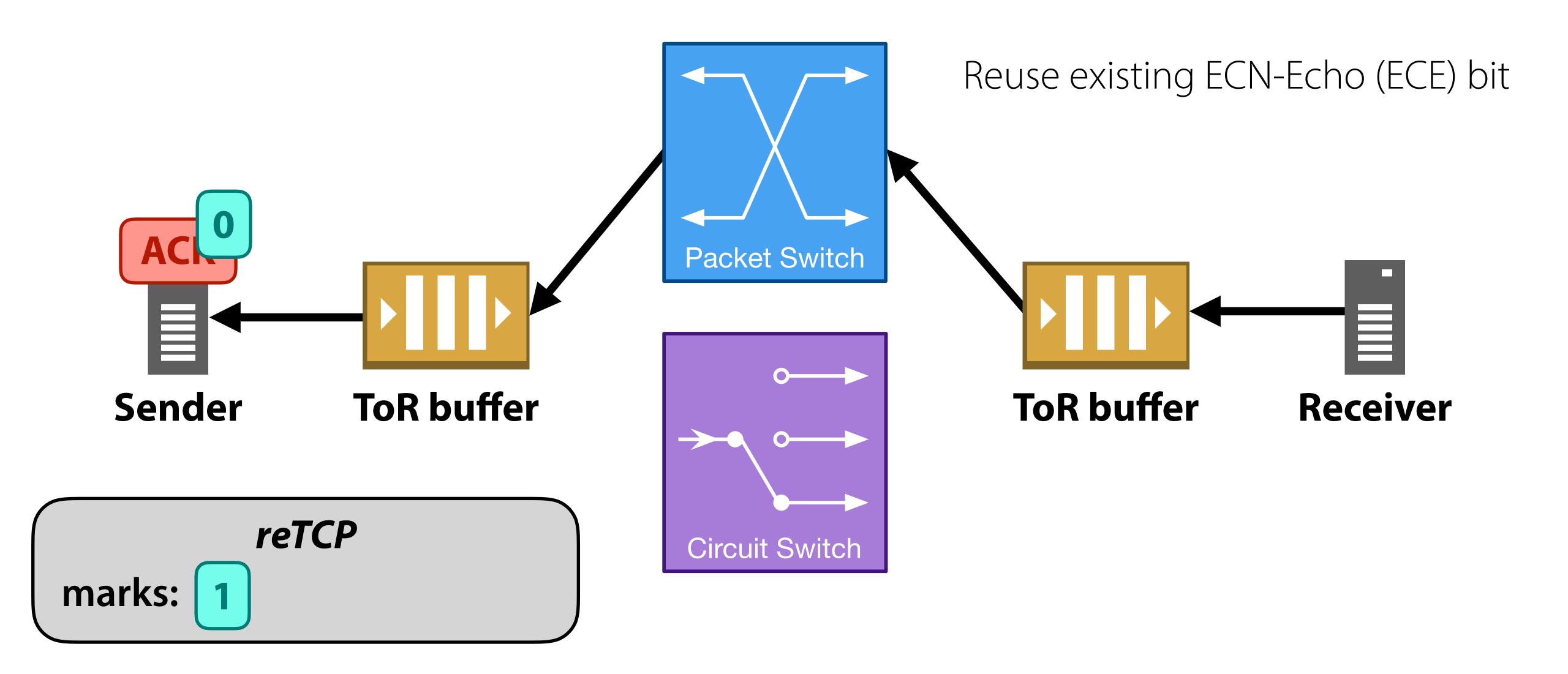


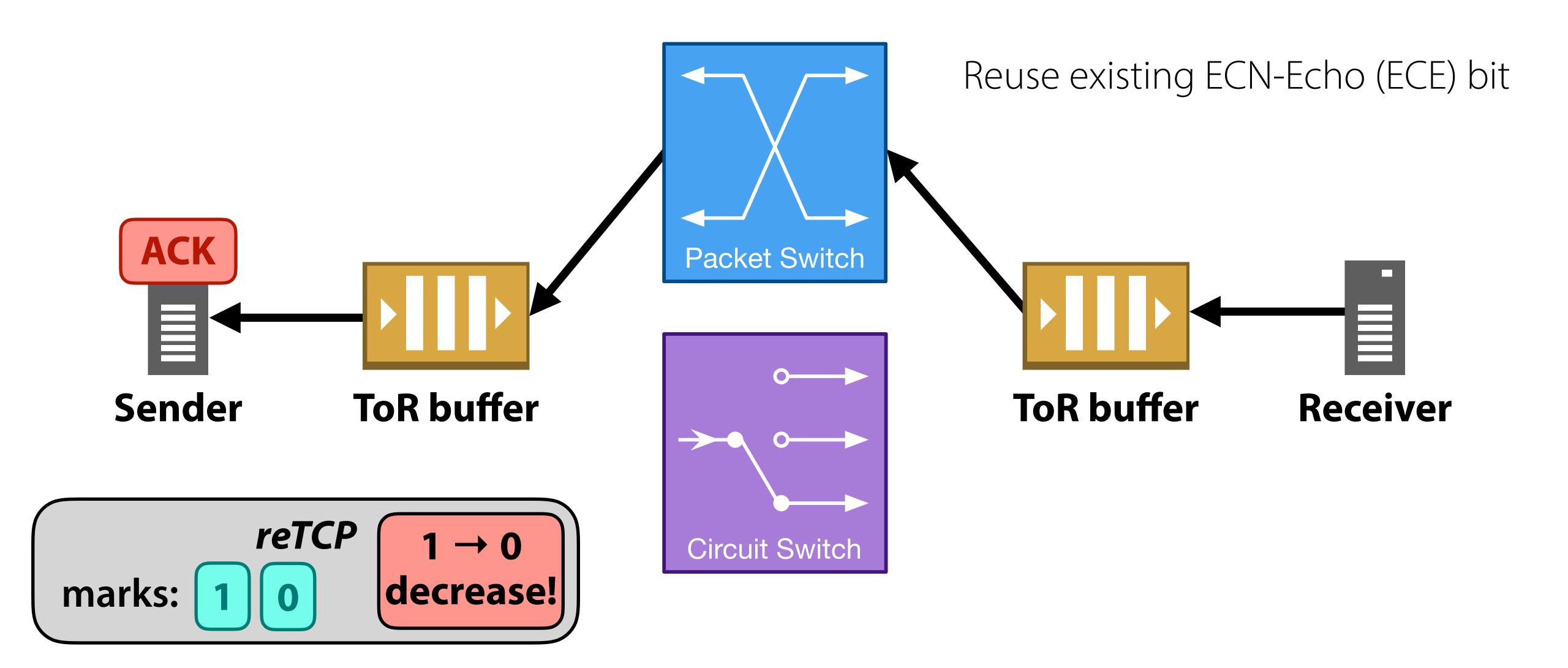












Single multiplicative increase/decrease

```
On 0 \rightarrow 1 transitions: On 1 \rightarrow 0 transitions: cwnd = cwnd × \alpha cwnd = cwnd / \alpha
```

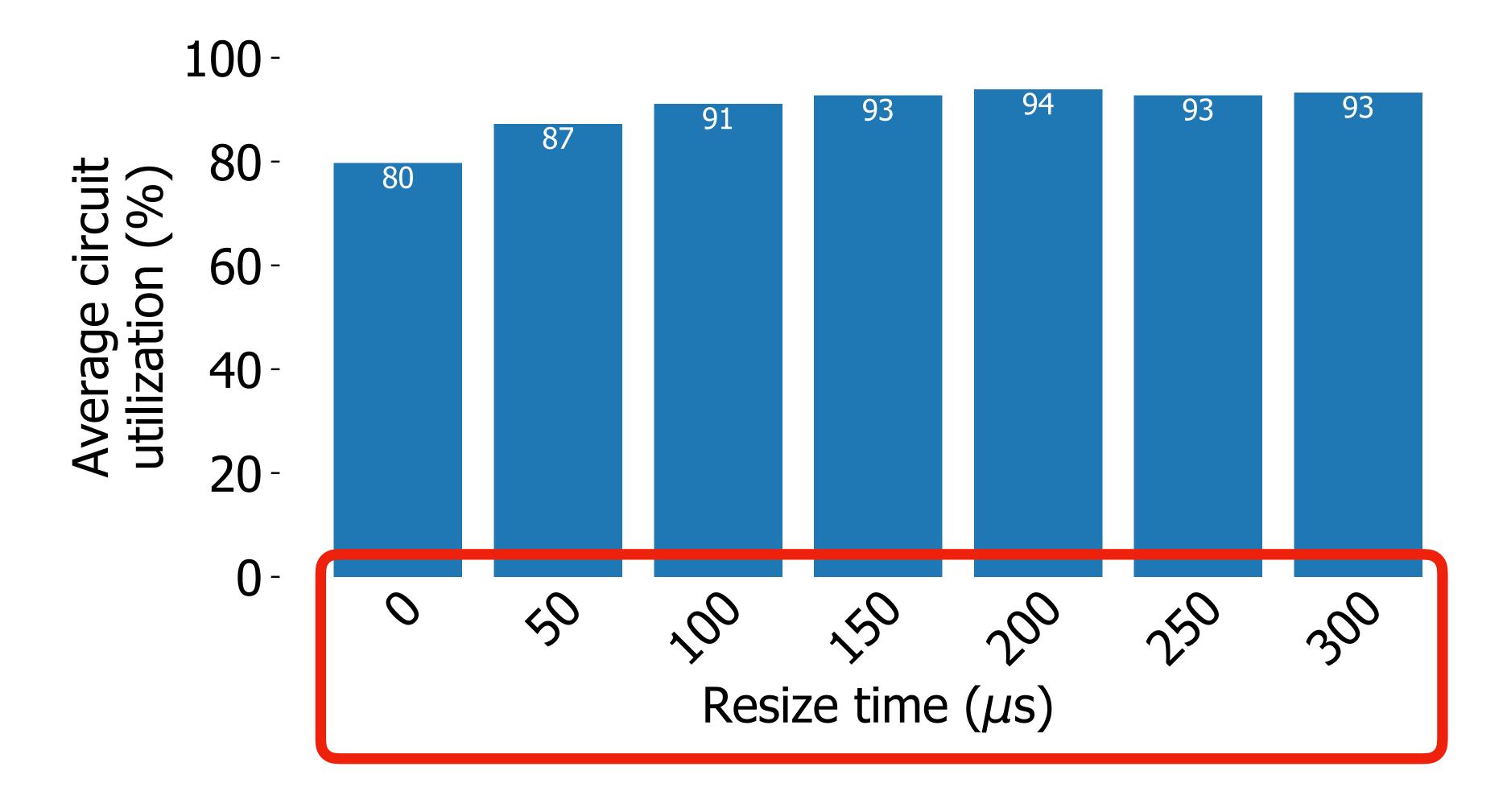
 α depends on ratio of circuit BDP to ToR queue capacity:

- Circuit network BDP: 45 packets
- Small ToR queue capacity: 16 packets

We use $\alpha = 2$

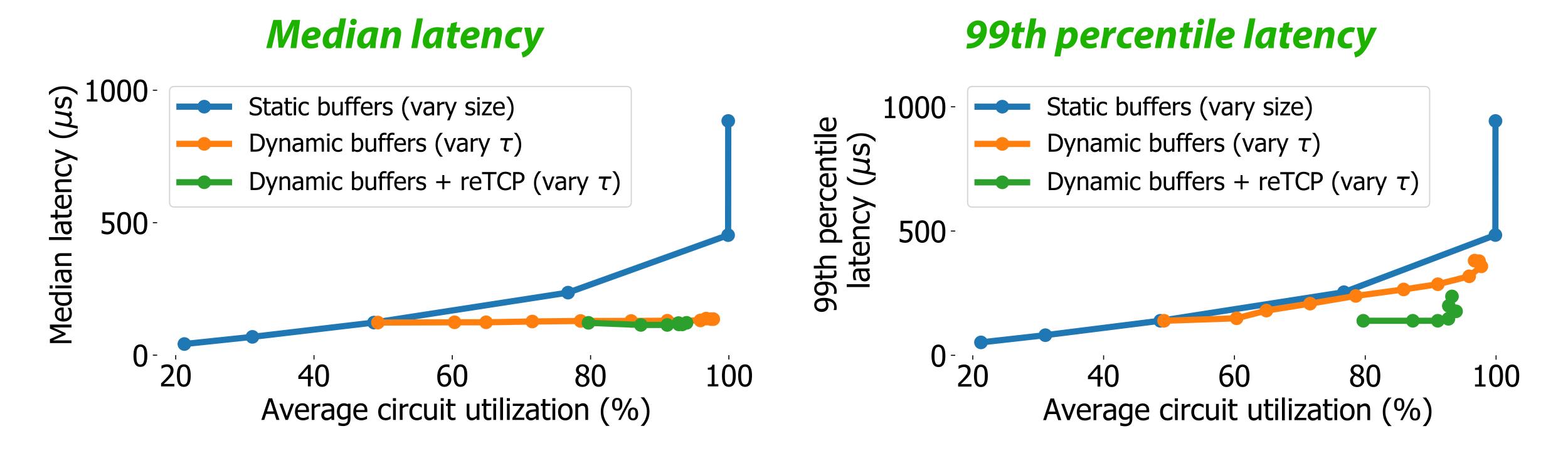
More advanced forms of feedback are possible

Dynamic buffers + reTCP achieve high utilization



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Short prebuffer time means low latency



With 150µs of prebuffering, dynamic buffers + reTCP achieve 93% circuit utilization with an only 1.20x increase in tail latency

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Limitations and future work

Dynamic buffer resizing and reTCP are designed to be minimally invasive

Higher performance may be possible by involving the end-host further

Our evaluation used a simple traffic pattern to isolate TCP's behavior

• Important to consider complex workloads as well

Is TCP the right protocol for hybrid networks?

Summary: Adapting TCP for RDCNs

Bandwidth fluctuations in reconfigurable datacenter networks break TCP's implicit assumption of relative network stability

Two techniques to ramp up TCP during short-lived circuits

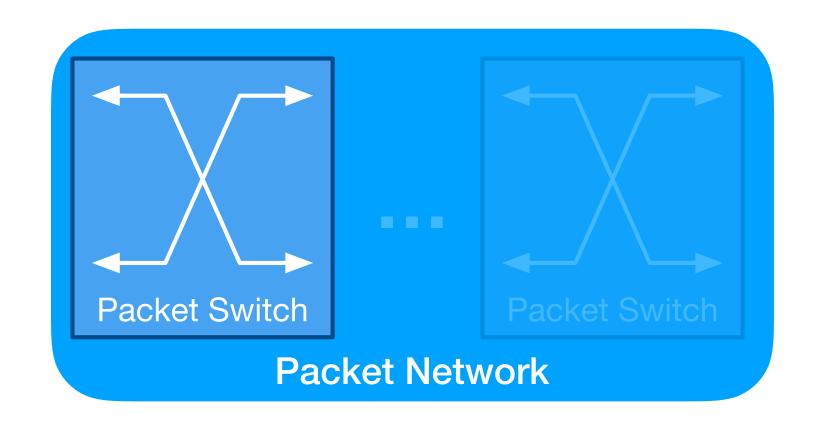
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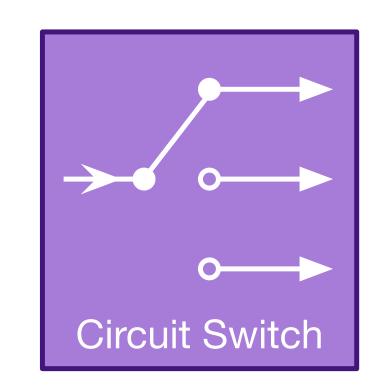
Etalon emulator open source at: github.com/mukerjee/etalon

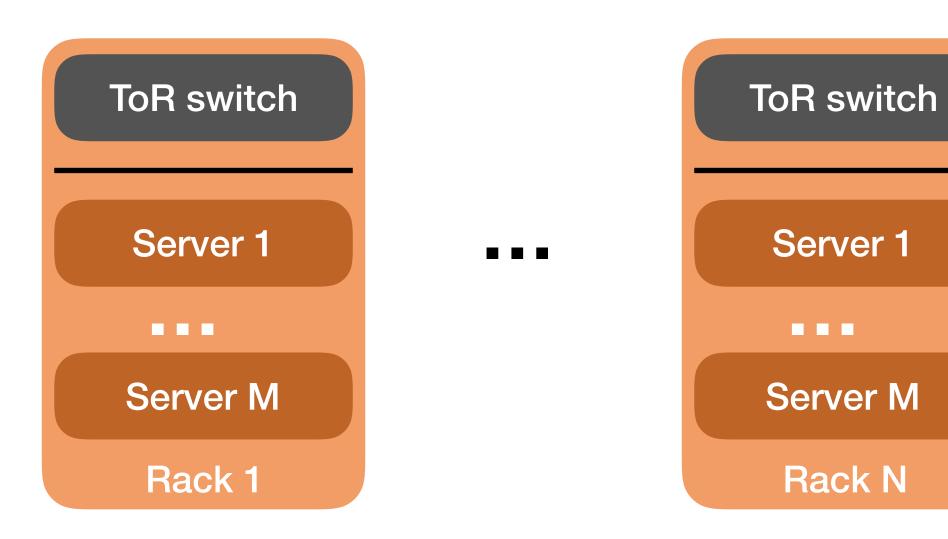
Christopher Canel ~ ccanel@cmu.edu

Thank you!

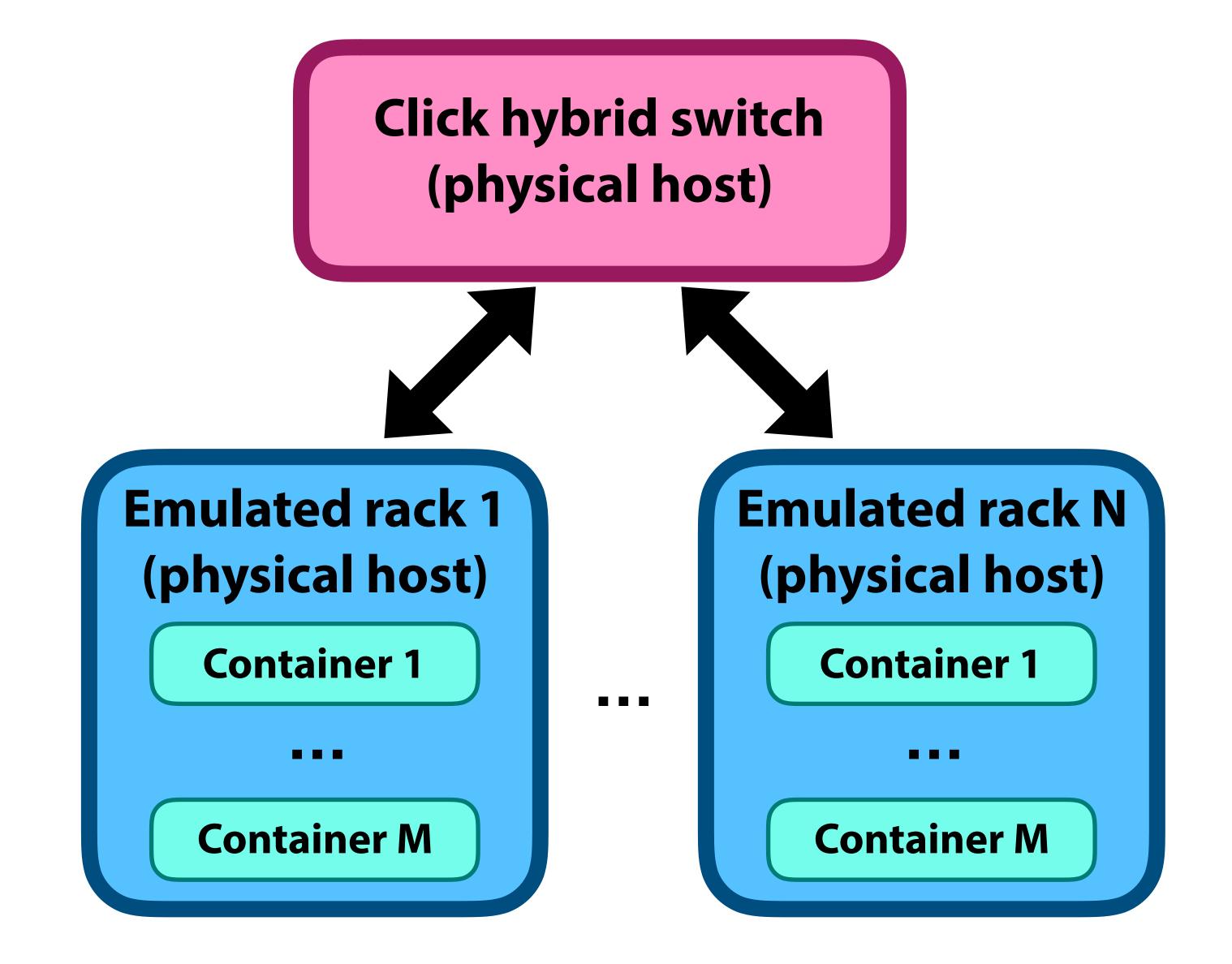
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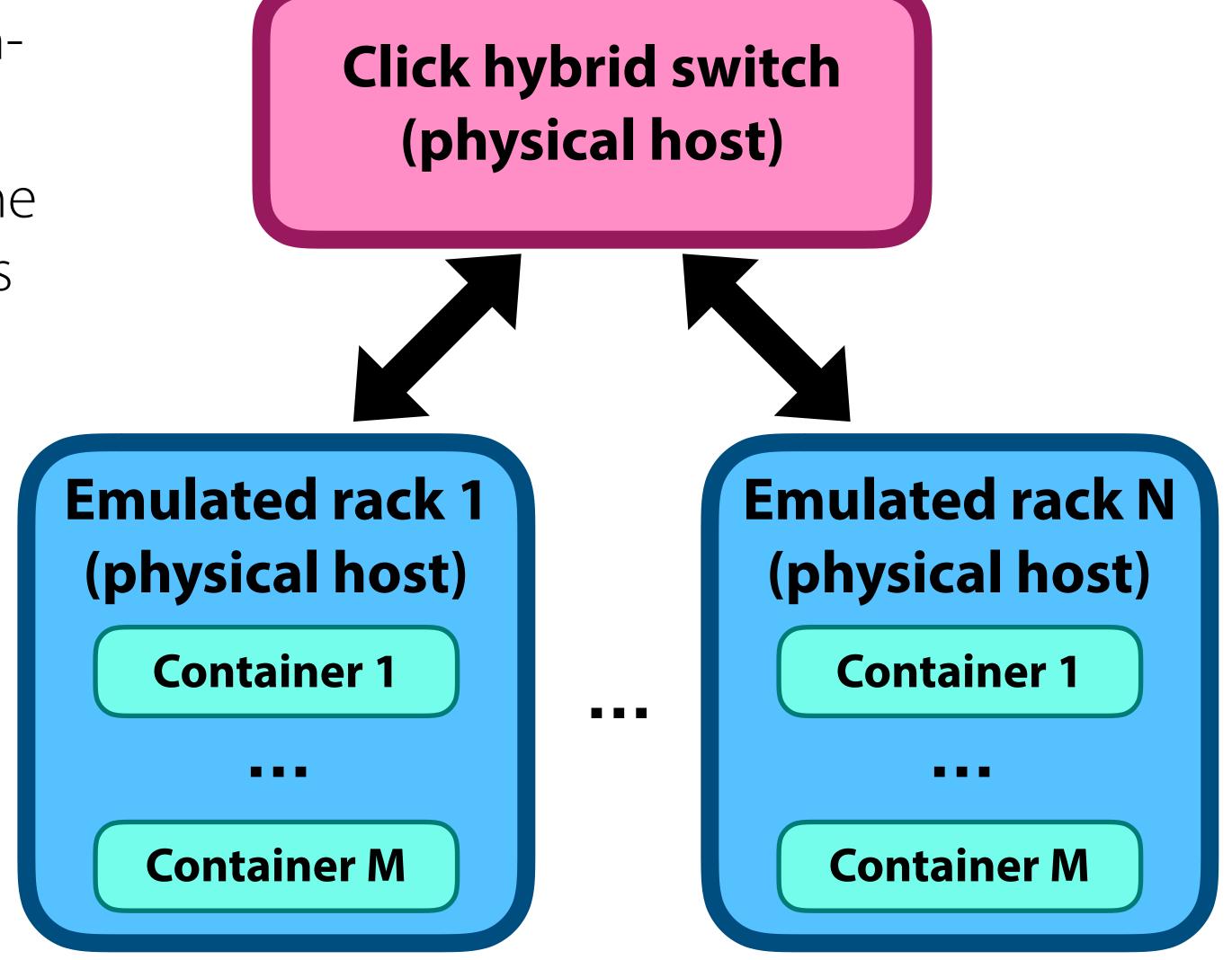
One more thing: Etalon emulator

Use *time dilation* to emulate high-bandwidth links

- "slows down" rest of the machine
- libVT: Catches common syscalls

Flowgrind to generate traffic

Strobe schedule: Each rack pair gets a circuit for an equal share



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